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A New Adaptive Clustering Technique for Large-Scale Sensor Networks

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#### **Outline of The Presentation**

Introduction

- System Architecture
- Adaptive Clustering
  - Real-Time Estimation of Number of Nodes
  - The Adaptive Clustering Algorithm

Simulation Results

- Example Network 1
- Example Network 2

Conclusion



Problems in Networking Sensor Systems:

- A large numbers of small, relatively inexpensive and low-power sensors are connected via a wireless network.
- Through the network, the data extracted from these sensors is sent to a nearby base station (BS), which forwards the data to a remote data center for further processing.

Challenging Problem: how to dynamically organize these sensors into a wireless communication network and route detected event information between field sensors and BS, i.e., the design of a networking protocol.

## **Networking Requirements**

Primary Requirements: energy efficiency, scalability, adaptivity, and security.

- Secondary Requirements: channel efficiency and network performance, which includes packet delay, packet loss ratio, throughput, and fairness.
- The focuses of this paper: energy efficiency, scalability, and adaptivity.

• Energy-efficient: the energy consumption in delivering packets from source to destination is minimized.

• Power-Aware: a route with nodes having higher residual (battery) power should be selected, although it may not be the shortest route.

## **Current Work**

Clustering is an energy-efficient architecture, wherein the individual nodes forward the information to their respective cluster heads.
The information is aggregated at the cluster head and then sent to the BS by the cluster head.

• The cluster heads and BS usually form a multi-hop network, which must have a multi-hop routing protocol.

## **Current Work: Clustering**

- The near term digital radio (NTDR): based on a two-tier clustering without using low-energy routing or MAC.
- *d*-hop clusters: directly aimed at minimizing the energy spent in the system.
- **LEACH:** the nodes were organized into cluster hierarchies and TDMA was applied within each cluster. Assumed the number of nodes and the optimal number of clusters to be formed are given.
- HEED: an improvement based on LEACH.
   Clustering uses node degree cost, which needs to know the number of active nodes. Not guarantee optimal selection of cluster heads in terms of energy.

The significant differences between the two networks:

• Most traffic in sensor network is triggered by sensed events;

• In reporting phase, the traffic goes from a hot-spot area, which consists of a few of the sensors, to a BS;

• In polling phase, traffic goes from a BS to many sensors;

• Many sensors coordinate for a common task and not all of them must be active.

Therefore, a sensor network is more like the fault management system (FMS) of a data network rather than the data network itself.

#### **Sensor and Data Networks**

- Many of the existing sensor networking protocols have often applied data network concepts and produced inefficient algorithms.
- In a typical application of a sensor network, the routed event can be either an anomaly event autonomously sent out by sensors, or a polling command issued by the remote collection stations.
- We propose the ACT and port many FMS concepts and protocols to sensor networks, such as event routing, reporting, processing, correlation, by using SNMP traps and alarms.



Two-tier hierarchical architecture, includes an adaptive clustering technique and multi-hop routing protocol.

- ACT: adaptive clustering technique.
- In tier-1, cluster members communicate to each other over a hybrid MAC protocol consisting of both contention and schduling.
- In tier-2, cluster heads communicate over a multi-hop routing protocol.
- Focus: an adaptive clustering technique for large-scale sensor networks.

## **System Architecture**



#### Figure 1: Architecture of A Sensor Network



## Assumptions

Network area: a square area of side 2a, the area  $A = 4a^2$ .

- The sensors inside the area are distributed according to a homogeneous spatial Poisson process, with intensity of  $\lambda \ sensors/m^2$ , hence, the mean value of n is  $\lambda A$ .
- The base station (BS) is at the center of the square area.

### Notations

- n: the number of sensors in the network, a random variable. For a particular realization, there are N sensors.
- *p*: the probability of a node becoming cluster head. Thus, on average, there are Np nodes that will become cluster heads, the rest N(1 − p) nodes will be cluster members in total.
- k: the average number of cluster heads. Therefore, there are k = Np clusters, each on average has m = N/k 1 member nodes; m is the average number of CMs within a cluster;

# **Adaptive Clustering**

The critical problem is how to estimate the parameters in real-time, N, the number of active nodes, and the optimal number of clusters,  $k_{opt}$ , or equivalently  $p_{opt}$ .

Assume that each node

• is capable of measuring its received signal power.

• keeps a record of the minimum power of the signals it has received within its radio range during previous cluster updating cycle.

## **Radio Propagation Model**

Notations:

•  $s_t$  and  $s_r$  are random variables that describe the powers of a signal a sensor node has transmitted and received at distance r, respectively;

•  $\kappa$  is a dimensionless constant which depends on the antenna characteristics and average attenuation from blockage, while  $r_0$  is a reference distance fro the antenna far-field;

•  $\gamma = 2 \sim 4$  is the path-loss exponent; for simplicity and without lossing generality, we choose  $\gamma = 2$ .

•  $\psi_{dB}$  is a Gaussian-distributed random variable with mean zero and variance  $\sigma_{\psi}$ , which can be also measured.

## **Radio Propagation Model**

The widely used measurement-based radio propagation model is the path-loss model with log-normal shadowing:

$$\frac{s_r}{s_t}(dB) = 10\log_{10}\kappa - 10\gamma\log_{10}\frac{r}{r_0} + \psi_{dB}, \quad (1)$$

By taking mean values on the random variables in above Eqn.,

$$\frac{S_r}{S_t}(dB) = 10\log_{10}\kappa - 10\gamma\log_{10}\frac{r}{r_0},$$
 (2)

where  $S_r$  and  $S_t$  are the mean values of  $s_r$  and  $s_t$ , respectively.

## **Radio Propagation Model**

Notations:

- S: the signal power a node has received.
- r: the distance of a node to its cluster head.
- Within a cluster, each node communicates to its cluster head by one hop distance; we have the further simplified model:

$$S(r) = \epsilon_0 / r^{\gamma}, \quad r \le r_{max}, \tag{3}$$

where

•  $\epsilon_0$ : a constant related to the radio trasmission device and propagation environments;

•  $r_{max}$ : the maximum allowable distance between a node and its cluster head.

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#### Voronoi Cell

Definition: a cell that covers most of the nodes in a cluster.

- $\rho_m$ : the radius of the minimal ball that covers the Voronoi cell. The ball is centered at the nucleus of the cell.
- The probability that  $\rho_m$  is greater than a certain value  $\rho$  has an upper bound:

$$Prob\{\rho_m > \rho\} \le 1 - [1 - \exp(-\mu p \lambda \rho^2)]^7,$$
 (4)

where

μ = 2(π/7 + sin π/14 + cos 5π/14),
pλ: the equivalent intensity for the point process that describes the cluster head nodes.

The above Eq. can be simplified:

$$Prob\{\rho_m > \rho\} \le 7\exp(-\mu p\lambda\rho^2).$$
 (5)

- Objective: to ensure that the probability of all the member nodes are beyond the minimal ball is very low,
- Define a parameter (degree of clustering)  $\alpha$  that:

$$Prob\{\rho_m > \rho\} \le \alpha,\tag{6}$$

where  $\alpha$  is a very small value specified from cluster design requirement.

By combining the two Eqns., we can choose

$$\alpha = 7 \exp(-\mu p \lambda \rho^2). \tag{7}$$

**We** can find such a  $\rho$  that must meet

$$\rho = \sqrt{-\ln(\alpha/7)/(\mu p \lambda)}.$$
(8)

In this way, we can make sure that  $\rho_m \leq \rho$  in most cases, i. e., the minimal ball covers most of the nodes in the Voronoi cell.

#### **Maximum Hop Distance**

Equivalently, the maximum distance of a node to a cluster head must be bounded by the above Eqn. in order to fall inside of the ball, that is,

$$r_{max} \le \sqrt{-\ln(\alpha/7)/(\mu p \lambda)}.$$
 (9)

Almost all the nodes are within one hop distance from a cluster head. Therefore, we can let  $r_{max}$  take the maximum value. Using the equal sign in above Eqn.,

$$S_{min} = -\frac{\epsilon_0 \mu p \lambda}{\ln(\alpha/7)},\tag{10}$$

where  $S_{min}$  is the minimum signal power a node has received.

Therefore, we can approximately estimate the value of the number of active nodes, by using  $N = \lambda A$ , if the probability of a node becoming cluster head, i.e., *p* value, is given.

## **Energy Efficient Clustering**

- In the network level, the requirement on clustering is to have as less number of clusters as possible.
- Within a cluster, the CMs must be as close to their CH as possible.
- To measure the cost incurred by the two requirements:

$$C(q) = e_1 k |D_0| + e_2 \sum_{j=1}^k m |D_1|, \qquad (11)$$

where  $D_0$  is the average distance from a CH to the BS; and  $D_1$  is the average distance of a CM to its CH.

## **Energy Efficient Clustering**

After some manipulations, we have

$$C(p) = e_1 d_0 anp + e_2 np(\frac{1}{p} - 1)\frac{2}{3}r_{max}.$$
 (12)

Now the energy efficient clustering becomes an optimization problem: how to choose an optimal p to minimize the total energy consumptions spent in intra- and inter-cluster communications?

#### **How to Choose Energy Coefficient?**

Assume that the average energy consumption in network layer for a CH to receive and transmit a unit of data (e. g., one bit, or one packet) over a hop distance of H are  $\overline{E}_r$  and  $\overline{E}_t$ , respectively.

**The total energy spent in network layer:** 

$$e_1 = (\overline{E}_r + \overline{E}_t)/H, \qquad (13)$$

where  $E_r = E_{elec}$ ,  $E_{elec}$  is the energy spent on electrical device for receiving a unit of data; and  $\overline{E}_t = E_{elec} + \epsilon_0 H^{\gamma}$ ,  $\epsilon_0$  and  $\gamma$  can take different values, depending on H, as defined in Table ??.

#### **How to Choose Energy Coefficient?**

Assume that the average energy consumption in MAC layer for a CM is  $\overline{E}_r$  for the receiving of the data bit per time unit (e. g., a time slot);  $\overline{E}_t$  for the successful transmission of the bit per time unit;  $\overline{E}_c$ for a collision that lasts for one time unit; and  $\overline{E}_i$  for an idle that lasts for one time unit.

The total energy spent in MAC layer:

 $e_2 = \left(w_s(\overline{E}_r + \overline{E}_t) + w_i\overline{E}_i + w_c\overline{E}_c\right)/\overline{D}_1, \quad (14)$ 

It can be found that

$$p_{opt} = 1/(\sqrt{c_1 S_{min}} - 1),$$
 (15)

where  $c_1 = -4c_0^2 A \ln(\alpha/7)/(\epsilon_0 \mu)$ .

Assume that during the *j*th cluster updating cycle, the measurement of  $S_{min}$  is denoted as  $\tilde{S}_{min}(j)$ , the corresponding value of  $p_{opt}$  as  $\tilde{p}_{opt}(j)$ .

$$\tilde{N}(j) = -\frac{\ln(\alpha/7)A\tilde{S}_{min}(j)}{\epsilon_0\mu\tilde{p}_{opt}(j)},$$
(16)

where N(j) is the calculated value of N in the jth cluster updating cycle.

By using a moving averaging model:

$$\hat{N}(j+1) = \beta \hat{N}(j) + (1-\beta)\tilde{N}(j+1), \quad (17)$$

where  $0 < \beta < 1$  is a smoothing factor used to adjust the estimation speed and accuracy.

The optimal number of clusters can be easily obtained

$$\hat{k}_{opt}(j+1) = \hat{N}(j+1)\tilde{p}_{opt}(j+1).$$
 (18)

Note that we can also have

$$\tilde{p}_{opt}(j+1) = 1/(\sqrt{c_1 \tilde{S}_{min}(j)} - 1).$$
 (19)

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# **The Adaptive Clustering Algorithm**

- 1. Initially, specify the value of  $\alpha$ , such as  $\alpha = 0.001$ . 2. Each node measures and records the minimum value  $S_{min}(j)$  in the current cycle.
- 3. Each node computes its  $p_{opt}$  value.
- 4. Each node computes its N and  $k_{opt}$  values.
- 5. By checking the inequality:

$$|\hat{N}(j) - \hat{N}(j-1)| \le \delta.$$
(20)

where  $\delta$  is a predefined QoS parameter. If yes, don't activate a cluster updating process. Go to step 2. 6. The node adopts the  $p_{opt}(j)$  value and tries to become a cluster head with this new probability. 7. Let j = j + 1, go to step 2.

#### Remarks

- Each node actively monitors its received signal power levels during the cluster updating cycles.
- The updating process is triggered by some serious changes in network status, by trading off energy consumption and the adaptivity.
- Each node attempts to find and stay at its optimal cluster locally. The algorithm is carried out by all nodes simultaneously and independently.
- If no serious changes, the algorithm is aborted silently. Otherwise, the node will automatically adjust its cluster to adapt to the changes in real-time.

## Conclusion

We propose

• a networking architecture for large-scale sensor networks

• an adaptive clustering technique

The simulation results have demonstrated that the two-tier hierarchical architecture

- adapts to changes in network topology,
- scales well to large network sizes,

• is power-aware and energy-efficient.

Future work would be further investigating the applicability of the proposed technique and algorithm to a larger number of other sensor networks.



Figure 2: Optimal Probability and Total Energy Spent

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Figure 4: Optimal Probability and Number of Nodes

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Figure 5: Numbers of Clusterheads and Nodes

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Figure 6: Ratio of Clustering Energy to Total Energy

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Figure 7: Network Lifetime Until First Node Inactive

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Figure 8: Network Lifetime Until Last Node Inactive

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Figure 9: Adaptive Clustering





Figure 10: Energy-Efficient Routing Scenario 1

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Figure 11: Energy-Efficient Routing Scenario 2

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Figure 12: Energy-Efficient Routing Scenario 3

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Figure 13: Energy-Efficient Routing Scenario 4

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