On the Interpretation of Energy as the Rate of Quantum Computation

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Abstract Over the last few decades, developments in the physical limits of computing and quantum computing have increasingly taught us that it can be helpful to think about physics itself in computational terms. For example, work over the last decade has shown that the energy of a quantum system limits the rate at which it can perform significant computational operations, and suggests that we might validly interpret energy as in fact being the speed at which a physical system is "computing," in some appropriate sense of the word. In this paper, we explore the precise nature of this connection. Elementary results in quantum theory show that the Hamiltonian energy of any quantum system corresponds exactly to the angular velocity of state-vector rotation (defined in a certain natural way) in Hilbert space, and also to the rate at which the state-vector's components (in any basis) sweep out area in the complex plane. The total angle traversed (or area swept out) corresponds to the action of the Hamiltonian operator along the trajectory, and we can also consider it to be a measure of the "amount of computational effort exerted" by the system, or effort for short. For any specific quantum or classical computational operation, we can (at least in principle) calculate its *difficulty*, defined as the minimum effort required to perform that operation on a worst-case input state, and this in turn determines the minimum time required for quantum systems to carry out that operation on worst-case input states of a given energy. As examples, we calculate the difficulty of some basic 1-bit and n-bit quantum and classical operations in an simple unconstrained scenario.

Key words Time evolution operator, Margolus-Levitin theorem, Hamiltonian energy, action of the Hamiltonian operator, quantum logic gates, energy as computing, physics as computation, geometric phase, quantum computational complexity

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1 Introduction

Over the years, the quest to characterize the fundamental physical limits of information processing has also helped to give us a deeper understanding of physics itself. For example, Shannon's studies of the limits of communication [1] taught us that the entropy of a system can also be considered to be a measure of the expected amount of unknown or incompressible information that is encoded in the state of that system. Landauer's [2] and Bennett's [3] analyses of the lower limit to the energy dissipation of computational operations led to Bennett's resolution [4] of the famous Maxwell's demon paradox, via the realization that the demon's record of its past perceptions is a form of physical entropy, which must be returned to the environment when that information is erased. More recently, Margolus and Levitin [5] showed that the energy of a quantum system limits the rate at which it can perform computational "operations" of a certain type, namely, transitions between distinguishable (orthogonal) quantum states. In the last few years, several articles by Lloyd and colleagues [6-8] have elaborated on this theme by suggesting that we can think of all variety of physical systems (ranging from particles and black holes to the entire universe) as comprising natural computers, with each system's "memory capacity" given by its maximum entropy, and its "computational performance" given by its total energy. We should also note that Ed Fredkin has been promoting a universe-as-computer philosophy for many decades.

The concept of interpreting physics as computing is certainly an exciting theme to pursue, due to its promise of conceptual unification, but we would like to proceed carefully with this program, and take the time to understand the details of this potential unification thoroughly and rigorously. While taking care to get all of the details exactly right, we would like not only to establish that a given physical quantity "limits" or "relates to" a given informational or computational quantity, but also justify the even stronger statement that the physical quantity actually *is*, at root, a fundamentally informational or computational quantity, one that has been traditionally expressed in terms of operationally defined physical units for reasons that can be viewed as being merely historical in nature.

As one the most famous examples of this type of conceptual progression, Rudolph Clausius [9] first defined (differential) entropy as the ratio of differential heat to temperature, dS = dQ/T, and at the time, entropy had no further explanation. Later, Ludwig Boltzmann [10] proposed the relation $S \propto -H = \int f \log f \, d\xi$ (where f is a probability density function ranging over particle energies or velocity vectors $\boldsymbol{\xi}$), which was backed up by his "H-theorem" showing that H spontaneously decreases over time for statistical reasons. In subsequent decades, this relation for entropy evolved and was generalized to become Boltzmann's eventual epitaph $S = k \log W$, which related entropy to the logarithm of the number of ways W of arranging a system [11].¹ Boltzmann's logarithmic quantity H (in a discrete and negated form) was later recognized by Shannon and others to also be an appropriate measure of the information content of a system. But, Boltzmann's fundamental insight regarding the nature of entropy can be viewed as having gone far beyond just *relating* a physical quantity to an information-based one. Rather, it can be viewed as telling us that physical entropy, at root, is really nothing but an informational quantity, one which merely manifests itself in terms of measurable physical units of heat and temperature due to the fact that these quantities themselves have an origin that is ultimately of a statistical nature, e.g., heat as disorganized energy.

Indeed, the long-term quest of physics to eventually create a grand unified "theory of everything" can be viewed as the effort to eventually reveal all physical concepts, quantities, and phenomena as being manifestations of underlying structures and processes that are purely mathematical and/or statistical in nature, and that therefore have an informational/computational flavor, at least insofar as the entire realm of formal mathematics can be viewed as being a fundamentally "computational" entity. As one interesting logical conclusion of this conceptual progression, if all observed phenomena are indeed eventually explicable as being aspects of some underlying purely mathematical/computational system, then we can argue that in the end, there really is no need for a separate physical ontology at all any more; we could instead validly suppose that the entire "physical" world really is nothing but a certain (very elaborate and complex) abstract mathematical or computational object. Such a viewpoint has many attractive philosophical features, at least from the perspective of a hard-core rationalist. One prominent proponent of such musings is Tegmark, e.g., see [12]. Another proposal for unifying mathematics and physics was recently made by Benioff [13].

However, regardless of one's personal feelings about such far-ranging philosophical agendas, if we can at least show that it is consistent to say that a given

¹ The references to Clausius and Boltzmann in this paragraph are also taken from [11].

physical quantity can be exactly identified with a given mathematical or computational quantity, then, as scientists, we can certainly all agree that the most parsimonious description of physics will indeed be one that does make that identification, since otherwise our description of the world would be burdened with an unnecessary proliferation of artificially distinct concepts, in violation of Ockham's razor, the most fundamental principle of scientific thought.

In this paper, we will primarily concern ourselves with just one small aspect of the grander theme of interpreting physics as information processing. Specifically, we focus on the idea of interpreting the physical energy content of a given system as being simply a measure of the rate at which that system is undergoing a certain ubiquitous physical process—namely, quantum state evolution—which can also be viewed as a computational process, as we do in quantum computing. In other words, the premise is that physical energy is nothing but the *rate of quantum computing*, if the meaning of this phrase is appropriately defined. This paper will clarify precisely in what sense this statement is true.

We'll also see that the concept of physical *action*, in a certain (somewhat generalized) sense, corresponds to a computational concept of the *amount of computational effort exerted*, which we'll call *effort* for short.

Of course, it is not necessarily the case that a given system will have been prepared in such a way that all of its physical computational activity will actually be directly applied towards the execution of a target application algorithm of interest. In most systems, only a small fraction of the system's energy will be engaged in carrying out application logic on computational degrees of freedom, while the rest will be devoted to various auxiliary supporting purposes, such as maintaining the stability of the machine's structure, dissipating excess heat to the environment, *etc.*, or it may simply be wasted in some purposeless activity.

For that part of energy that *is* directly engaged in carrying out desired logical operations, we will see that one fruitful application of the computational interpretation of energy will be in allowing us to characterize the *minimum* energy that must be harnessed in order to carry out a given computational operation in a given period of time. In section 12, we will show how to calculate this "difficulty" figure for a variety of simple quantum logic operations, and we briefly discuss how to generalize it to apply to classical reversible and irreversible Boolean operations as well.

2 Background

Of course, the earliest hints about the relationship between energy and the rate of computing can be found in Planck's original $E = h\nu$ relation for light, which tells us that an electromagnetic field oscillation having a frequency of ν requires an energy at least $h\nu$, where $h \simeq 6.626 \times 10^{-34}$ J s is Planck's constant. Alternatively, a unit of energy E, when devoted to a single photonic quantum, results in an oscillation (which can be considered to be a very simple kind of computational process) occurring at a cycle rate of $\nu = E/h$.

Also suggestive is the Heisenberg energy-time uncertainty principle $\Delta E \Delta t \ge h/2$, which relates the standard deviation or uncertainty in energy ΔE to the min-

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imum time interval Δt required to measure energy with that precision; the measurement process can be considered a type of computation. However, this relation by itself only suggests that the *spread* or standard deviation of energy has something to do with the rate of a process of interest; whereas we are also interested in finding a computational meaning for the absolute or mean value of the energy, itself.

More recently, in 1992, Tyagi [14] proposed a notion of "computational action" that was based on the amount of energy *dissipated* multiplied by the elapsed time (a quantity which has the same physical units as action) and proposed a theory of optimal algorithm design based on a "principle of least computational action." However, Tyagi's analogy with Hamilton's principle was still a long way from indicating that *physical* action actually *is* computation in some sense, or that physical energy itself (which is, in general, not necessarily dissipated) corresponds to a rate of computation. Still, it was suggestive.

Going much further, in 1998 Toffoli [15] argued that the least-action principle in physics itself can be derived mathematically from *first principles* (rather than as an *ad hoc* physical postulate) as a simple combinatorial consequence of counting the number of possible fine-grained discrete dynamical laws that are consistent with a given macroscopic trajectory. In Toffoli's model, which intriguingly even captures aspects of relativistic behavior, the energy of a state is conjectured to represent the logarithm of the length of its dynamical orbit. Toffoli also gives a correspondence between physical action and amount of computation that is more explicit than Tyagi's, and in which the path with the least Lagrangian action is the one with the greatest amount of "unused" or "wasted" computational capacity. In later papers following up on the present one, we will show that indeed, Lagrangian action corresponds negatively to the portion of the computational effort that does not contribute to an object's active motion.

At around the same time as Toffoli's work, Margolus and Levitin [5] showed that in any quantum system, a state with a quantum-average energy E above the ground state of the system takes at least time $\Delta t \ge t^- = h/4E$ to evolve to an orthogonal state, along with a tighter bound of $\Delta t \ge t^-_N = (N-1)h/2NE$ that is applicable to a trajectory that passes through a cycle of N mutually orthogonal states before returning to the initial state. In the limit as $N \to \infty$, $t^-_N \to h/2E$, twice the minimum time of $t^- = t^-_2$ which applies to a cycle between 2 states. Both bounds are achievable in principle, in freely constructed quantum systems.

In a widely-publicized paper in *Nature* in 2000, Lloyd [6] used the Margolus-Levitin result to calculate the maximum performance of a 1 kg "ultimate laptop," in a hypothetical limiting scenario in which all of the machine's rest mass-energy is devoted to carrying out a desired computation.

Two years later, Levitin, Toffoli and Walton [16] investigated the minimum time to perform a specific quantum logic operation, namely a CNOT (controlled-NOT) together with an arbitrary phase rotation, in systems of a given energy E.

In 2003, Giovannetti, Lloyd and Maccone [17,18] explored tighter limits on the time required to reduce the fidelity between initial and final states to a given level, taking into account the magnitudes of both E and ΔE , the system's degree of entanglement, and the number of interaction terms in the system's Hamiltonian. Results such as the above suggest that energy might fruitfully be *exactly* identified with the rate of raw, low-level quantum-physical "computing" that is taking place within a given physical system, in some appropriate sense, if only the quantity "amount of computing" could be defined accordingly. We would like to show that some well-defined and well-justified measure of the rate at which "computational effort" (not necessarily useful) is being exerted within any quantum system is indeed *exactly* equal to the energy of that system.

3 Preview

In subsequent sections of this paper, we address the aforementioned goal by proposing a well-defined, real-valued measure of the total *amount of change* undergone over the course of *any* continuous trajectory of a normalized state vector along the unit sphere in Hilbert space. This measure is simply given by the line integral of the magnitude of the imaginary component of the inner product between infinitesimally adjacent normalized state vectors along the given path. This quantity is invariant under any time-independent change of basis, since the inner product itself is. As we will show, it is also numerically equal to twice the complex-plane area (relative to the origin) that is circumscribed or "swept out" by the coefficients of the basis vector components, in any basis. For closed paths, this quantity is even invariant under not only rotations but also translations of the complex plane. Finally, our quantity can be perhaps most simply characterized as being *the action of the Hamiltonian* along the path; this is to be contrasted with the usual action (of the Lagrangian), whose precise computational meaning will be addressed in later work.

We propose that the above-described measure of "amount of change" is the most natural measure of the amount of computational *effort* exerted by a physical system as it undergoes a specific trajectory. For any pair of trajectory endpoints, the effort has a well-defined minimum value over possible trajectories which is obtained along a "geodesic" trajectory between the endpoint states, thereby inducing a natural metric over the Hilbert space.

We will show that in any quantum system, the instantaneous rate at which change occurs (computational effort is exerted) for any state, under any time-dependent Hamiltonian operator, is exactly given by the (Hamiltonian) instantaneous average energy of the state. Thus, the state's energy *is* exactly its rate of computation, in this sense.

We use the word "effort" here rather than "work" both (a) to distinguish our concept from the usual technical meaning of work in physics as being directed energy, and also (b) to connote that effort is something that can be ineffectually wasted; *i.e.*, it does not necessarily correspond to *useful* computational work performed. In fact, we will see that indefinitely large amounts of effort could be expended (inefficiently) in carrying out any given quantum computational task, *i.e.* in accomplishing a given piece of computational work.

Despite having no upper bound, our concept of effort turns out to still be meaningful and useful for characterizing computational tasks, since (as we will see) any given quantum or classical computational operation does have a well-defined and non-trivial *minimum* required effort for worst-case inputs, which we will call the *difficulty* of the operation. As we will see, for any pair of unitaries U_1, U_2 , the difficulty of the operation $U_2U_1^{\dagger}$ that takes us from U_1 to U_2 gives a natural distance metric over U_n , the Lie group of rank-*n* unitary operators.

The difficulty of a computational operation, according to our definitions, determines the minimum time required to perform it on worst-case inputs of given energy, or (equivalently) the minimum worst-case energy that must be devoted to a system in order to perform the operation within a given time. The difficulty thus directly characterizes the computational complexity or "cost" of a given operation, in the same "energy-delay product" units that are popular in electrical engineering, but where the energy here refers to the average instantaneous energy that is *invested* in carrying out the computation, rather than to the amount of energy that is *dissipated*.

4 A Simple Example

In this section, we start by presenting a simple, concrete example in order to help motivate our later, more general definitions. Consider any quantum system subject to a constant (time-independent) Hamiltonian operator H. Let $|G\rangle$ and $|E\rangle$ be any normalized, non-degenerate pair of the system's energy eigenstates. The labels G and E here are meant to suggest the ground and excited states of a non-degenerate two-state system, but actually it is not necessary for purposes of this example that there be no additional states of higher, lower, or equal energy.

Since the Hamiltonian is only physically meaningful up to an additive constant, let us adjust the eigenvalue corresponding to vector $|G\rangle$ to have value 0 (*i.e.* let $H|G\rangle = 0$), and then let E denote the eigenvalue of $|E\rangle$ (*i.e.*, $H|E\rangle = E|E\rangle$). For example, for a two-state system, we could let $H = (1 + \sigma_z)E/2$ with the usual definition of the Pauli z-axis spin operator $\sigma_z = \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix}$; and let $|G\rangle = \begin{bmatrix} 0 \\ 1 \end{bmatrix}$ and

 $|E\rangle = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$, thus we have that $H = |E\rangle\langle E|$ and so E = 1.

Now, consider the initial state $|\psi_0\rangle = (|\mathbf{G}\rangle + |\mathbf{E}\rangle)/\sqrt{2}$ at time t = 0, and let it evolve over time under the influence of the system's Hamiltonian, with $|\psi(t)\rangle = e^{iHt/\hbar}|\psi_0\rangle$ denoting the state vector at time t.² Let $c_{|\mathbf{G}\rangle}(t)$ and $c_{|\mathbf{E}\rangle}(t)$ denote $\langle \mathbf{G}|\psi(t)\rangle$ and $\langle \mathbf{E}|\psi(t)\rangle$ respectively, *i.e.*, the components (complex coefficients) of the state vector $|\psi(t)\rangle$ when decomposed in an orthonormal basis that includes $|\mathbf{G}\rangle, |\mathbf{E}\rangle$ as basis vectors.

Initially, $c_{|G\rangle}(t) = c_{|E\rangle}(t) = 1/\sqrt{2}$. Over time, $c_{|E\rangle}$ phase-rotates in the complex plane in a circle about the origin, at an angular velocity of $\omega_{|E\rangle} = E/\hbar$. In time t = 2E/h, it rotates by a total angle of $\theta = \pi$. The area swept out by the line between $c_{|E\rangle}(t)$ and the origin is $a_{|E\rangle} = \frac{1}{2}\pi |c_{|E\rangle}|^2 = \pi/4$. This is the area of a semi-circular half-disc with radius $r_{|E\rangle} = |c_{|E\rangle}| = 1/\sqrt{2}$. Meanwhile, $c_{|G\rangle}(t)$ is

 $^{^{2}}$ For convenience, we use the opposite of the ordinary sign convention in the time-evolution operator.

Fig. 1 Under the Hamiltonian $H = E|E\rangle\langle E|$, starting from the initial state $|\psi_0\rangle = (|G\rangle + |E\rangle) \cdot 2^{-1/2}$, the complex coefficient $c_{|E\rangle} = \langle E|\psi(t)\rangle$ of $|E\rangle$ (the excited state) in the superposition sweeps out a half-circle in the complex plane with area $\pi/4$ in time t = 2E/h, while the ground-state coefficient $c_{|G\rangle}$ remains stationary.

stationary and sweeps out zero area. The total area swept out by both components is thus $a = \pi/4$. This evolution is depicted in fig. 1.

Does the area swept out by the complex components of the state vector depend on the choice of basis? We will answer this question in a much more general setting later, but for now, consider, for example, a new basis that includes basis vectors $|0\rangle$, $|1\rangle$ where $|0\rangle = (|G\rangle + |E\rangle)/\sqrt{2}$ and $|1\rangle = (|G\rangle - |E\rangle)/\sqrt{2}$. Consider the evolution again starting from the same initial state as before, $|\psi_0\rangle = |0\rangle$. Note that the final state after time t = 2E/h is $|1\rangle$. In the new basis, the coefficients $c_{|0\rangle}(t)$ and $c_{|1\rangle}(t)$ respectively trace out the upper and lower halves of a circle of radius 1/2 centered at the point 1/2 + i0. The total area swept out by both components (on lines between them and the origin) is the area of this circle, namely $a = \pi (1/2)^2 = \pi/4$. (See fig. 2.) Note that the total area in this new basis is still $\pi/4$.

At this point we may naturally ask, is the area the same in *any* fixed basis? Later we will show that the answer is yes; in general, the area swept out is independent of the basis for *any* trajectory of *any* initial state. The area swept out will be (proportional to) our proposed measure of the amount of computational effort exerted by a system in undergoing any specific state-vector trajectory.

Fig. 2 The evolution from fig. 1, re-plotted in the basis $|0\rangle = (|G\rangle + |E\rangle) \cdot 2^{-1/2}$, $|1\rangle = (|G\rangle + |E\rangle) \cdot 2^{-1/2}$. The coefficients of $|0\rangle$ and $|1\rangle$ together sweep out a full circle, but the total area swept out is still $\pi/4$.

5 General Framework

In this section we proceed to set forth the general mathematical definitions and notations to be used in the subsequent analysis.

5.1 Time-independent case

Let \mathcal{H} be any Hilbert space. Any linear, norm-conserving, invertible, continuous and time-independent dynamics on such a space must proceed via the application of a unitary time-evolution operator, expressible as

$$U = U(\Delta t) = e^{iA(\Delta t)} = e^{iH\Delta t}$$
(1)

where Δt is the length of a given time interval, $A(\Delta t) = H\Delta t$ maps the interval to an Hermitian operator A that is proportional to Δt , and H is an Hermitian operator with units of angular frequency. For any two times $t_1, t_2 \in \mathbb{R}$, and for any initial state vector $|\psi\rangle = |\psi(t_1)\rangle$ at time t_1 , the implied state at any other time t_2 is given by $|\psi(t_2)\rangle = U(\Delta t)|\psi(t_1)\rangle$, where $\Delta t = t_2 - t_1$. We will sometimes also write U and A as functions of the directed pair of times, written $t_1 \rightarrow t_2$. We will sometimes call the U and A operators "cumulative" when the interval Δt is not infinitesimal.

Note that in eq. (1) we are using the opposite of the usual (but arbitrary) negative-sign convention in the exponent; this is an inessential but convenient choice, in that later it will let us automatically associate positive energies with

positive (*i.e.*, counter-clockwise) phase velocities for the coefficients of state components.

For convenience, for any operator O and vector v, we will sometimes use the notation O[v] as an abbreviation for the expectation value $\langle v|O|v\rangle$.

Now, of course, the eigenvectors of U are also eigenvectors of A and H, so H's expectation value $H[\psi]$ for any initial vector $\psi(t_1) \in \mathcal{H}$ is preserved by the time-evolution $\psi(t_1) \rightarrow \psi(t_2)$. This conserved quantity (whose existence follows from time-independence even more generally, via Nöther's theorem) is called the *Hamiltonian energy* of the system. Although in our expressions it has the dimensions of angular velocity, this is the same as energy if we choose units where $\hbar = 1$, as is customary. Thus, H is called the Hamiltonian operator. We will call the operator $A = A(t_1 \rightarrow t_2)$ the *cumulative action of the Hamiltonian from time* t_1 to t_2 , where some of the qualifying phrases may be omitted for brevity. The reasons for the use of the word "action" will be discussed later.

For convenience in the subsequent discussions, we will often just set $t_1 = 0$ (without loss of generality) and write $U = U(t) = U(0 \rightarrow t) = e^{iHt}$. We refer to the complete operator-valued function $\lambda t.U(t)$ for all t values in some range (which usually includes t = 0, for which U(0) = I) as a *unitary trajectory* over that time interval. Also, for any t we write $A(t) :\equiv A(0 \rightarrow t)$ for the cumulative action from 0 to t.

Differentiating U(t) with respect to time and applying the result to an initial state $|\psi(0)\rangle$ then yields us Schrödinger's equation in various forms that we'll use,

$$\dot{U} = \frac{\mathrm{d}U(t)}{\mathrm{d}t} = \frac{\mathrm{d}}{\mathrm{d}t}\mathrm{e}^{\mathrm{i}Ht} = \mathrm{i}H\mathrm{e}^{\mathrm{i}Ht} = \mathrm{i}HU(t) \tag{2}$$

$$\frac{\mathrm{d}}{\mathrm{d}t}U(t)|\psi(0)\rangle = \mathrm{i}HU(t)|\psi(0)\rangle \tag{3}$$

$$\dot{\psi}\rangle = \frac{\mathrm{d}}{\mathrm{d}t}|\psi(t)\rangle = \mathrm{i}H|\psi(t)\rangle$$
(4)

$$\frac{\mathrm{d}}{\mathrm{d}t} = \mathrm{i}H,\tag{5}$$

where again, note that we are using $\hbar = 1$ and the opposite of the usual sign convention. Note also that we are able to differentiate e^{iHt} in eq. (2) because d/dt commutes with H, since H here is a constant.

5.2 Time-dependent case

The natural generalization of eq. (5) (the operator form of Schrödinger's equation) to a system with a time-dependent Hamiltonian H(t) is of course just

$$\frac{\mathrm{d}}{\mathrm{d}t} = \mathrm{i}H(t) \tag{6}$$

where now H(t) is permitted to vary over time, though often with a constraint that it be differentiable, smooth, or analytic.

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One may at first think that in this time-dependent context, we could appropriately generalize the time-evolution operator equation (1) by simply changing the definition of the action operator A (as a function of t) from the original A(t) = Htto what one might naïvely think would be the obvious generalization to a timedependent H,

$$A(t) = \int_{\tau=0}^{t} H(\tau) \mathrm{d}\tau, \tag{7}$$

while still keeping the relation $U(t) = e^{iA(t)}$. But in fact, the definition (7) does not work for this purpose, since in general the values of $H(\tau)$ at different times τ will not commute with each other; taking the integral loses all information about their relative time-ordering, and the time-derivative of U(t) will no longer be equal to iH(t) as required, since d/dt will no longer commute with H(t).

The standard way to repair this problem (discussed in almost any quantum field theory textbook, *e.g.*, [19]) is to define a time-ordering meta-operator \mathcal{T} , which takes a given operator expression and reorders its internal operator products so that operators associated with earlier time points are applied first in all products (reading right-to-left). For example, as a matter of definition,

$$\mathcal{T}[H(t_1)H(t_2)] := \begin{cases} H(t_1)H(t_2) \text{ if } t_1 > t_2\\ H(t_2)H(t_1) \text{ otherwise} \end{cases}$$
(8)

With this notational convention, we can write

$$U(t) = \mathcal{T} e^{iA(t)} \tag{9}$$

where A(t) is as defined in eq. (7), and the meaning of this meta-expression will be well-defined and consistent with eq. (6) applied to U(t). But the problem with this approach is that the expression A(t) in (9) no longer denotes a "first class object" of our language, but rather is a sort of meta-mathematical place-holder to be manipulated via a rather complex interpretational procedure, which involves applying eq. (8) to uncountably many infinitesimal pieces of the integrals appearing in the Taylor-expanded version of eq. (9). There is no longer any simple, direct relationship between the properties of the linear operator A(t) defined in eq. (7) (*e.g.*, its eigenvalues and eigenvectors) and the properties of U(t).

Thus, in what follows we will find it more useful to instead abandon eq. (7), and take the rather more concrete approach of simply redefining A(t) for a given unitary trajectory U(t) to be the unique continuously time-dependent Hermitian operator such that A(0) = 0 and

$$U(t) = e^{iA(t)} \tag{10}$$

(with *no* time-ordering operator!) for all t. To see that such an A indeed exists and is unique, note that since each particular U = U(t) (at a given moment) is unitary, it is a normal operator and can thus be given a spectral decomposition

$$U = \sum_{i} u_{i} |u_{i}\rangle \langle u_{i}| \tag{11}$$

where $\{|u_i\rangle\}$ and $\{u_i\}$ respectively comprise an orthonormal eigenbasis of U and the corresponding unit-modulus eigenvalues. We can therefore define the multi-valued logarithm of U by

$$\ln U = \ln \sum_{i} u_{i} |u_{i}\rangle \langle u_{i}|$$

$$:\equiv \sum_{i} (\ln u_{i}) |u_{i}\rangle \langle u_{i}|$$

$$= \sum_{i} i \arg(u_{i}) |u_{i}\rangle \langle u_{i}|$$
(12)

$$= \sum_{i} i[\operatorname{Arg}(u_i) + 2\pi n_i] |u_i\rangle \langle u_i|$$
(13)

where in step (12) we have used the fact that $|u_i| = 1$, and where in line (13) $\operatorname{Arg}(u_i) \in [0, 2\pi)$ denotes the principal value of the multivalued function $\operatorname{arg}(u_i)$, while the n_i values may be any integers. Although we see that there are infinitely many values of $(\ln U)$ for any individual U in isolation, nevertheless there is a unique single-valued definition of the entire function $L(t) = \ln U(t)$, given the function U(t), that is *continuous* over t and where L(0) = 0.

The uniqueness is due to the fact that U(t) varies continuously in t, and thus, if we like, the eigenbasis $\{|u_i(t)\rangle\}$ that we choose for U at each moment (which has k free gauge-like parameters determining the u_i , where $k = \dim \mathcal{H}$) can vary continuously as well. Given basis vectors $|u_i\rangle$ (and thus u_i values) that change continuously, it follows that at any moment, only one assignment of values to the n_i parameters can possibly yield continuity with the logarithm value L(t - dt)at the previous moment, since any other choice would (discontinuously) change one of the phase angles $\operatorname{Arg}(u_i) + 2\pi n_i$ in the expression (13) by an amount that is (infinitesimally close to) a multiple of 2π . The n_i parameters can (and must) change by ± 1 from their preceding values (while leaving L(t) continuous) only at a discrete set of time points, namely those where the continuously-changing u_i value crosses the branch cut of the Arg() function (in some direction), and $\operatorname{Arg}(u_i)$ jumps by $\mp 2\pi$.

Now, given this uniquely-defined unitary trajectory logarithm $L(t) = \ln U(t)$, we simply define our action operator as A(t) = -iL(t), and then trivially we have that $U(t) = e^{iA(t)}$ holds for all t, where the exponential can be defined via the spectral decomposition of A (equivalently to the standard Taylor-series definition), thereby inverting the logarithm.

Meanwhile, the entire unitary trajectory U(t) itself is derived from the Hamiltonian trajectory H(t) by setting U(0) = I and applying the operator form (6) of the time-dependent Schrödinger equation to U(t). So (d/dt)U(t) = iH(t)U(t), and we are thereby guaranteed that in fact

$$\frac{\mathrm{d}}{\mathrm{d}t}\mathrm{e}^{\mathrm{i}A(t)} = \mathrm{i}H(t)\mathrm{e}^{\mathrm{i}A(t)} \tag{14}$$

as desired, which (recall) failed to be true (in the absence of a time-ordering operator) for the A(t) defined in eq. (7). For reasons we will explain, we will refer to a complete function $\lambda t.A(t)$ as defined by eq. (10) as the *cumulative Hamiltonian action trajectory* implied by the Hamiltonian trajectory H(t).

In cases where H(t) = H is constant over time, note that this definition of A(t) reduces to the simple Ht form that we used back in eq. (1). This follows from the observation that the definition A(t) = Ht indeed solves eq. (10) when H is constant, and the fact that (as we just showed) the A(t) implied by eq. (10) is unique under the continuity constraint.

Later, we will see the importance of the Hamiltonian action trajectory A(t), and discuss the precise meaning and computational interpretation of its expectation value when applied to a given state.

To clarify our terminology, note that in this document we are using the word *action* in a somewhat more general sense than is usual; typically in physics (*e.g.*, in Hamilton's principle) "action" just refers to the quantity having units of action that is obtained by integrating the Lagrangian L = pv - H along some path. However, it is also perfectly valid and reasonable to consider the more general notion of the action that is associated with *any* quantity that has units of energy, by setting the time-derivative of that action along some path to be equal to that energy.

Indeed, we will see later that the time-derivative of the cumulative Hamiltonian action A(t) (as we have defined it) along a given trajectory is in fact exactly the instantaneous Hamiltonian energy H(t), *i.e.*,

$$\frac{\mathrm{d}}{\mathrm{d}t}A(t)[\psi(0)] = H(t)[\psi(t)],\tag{15}$$

similarly to how the time-derivative of the ordinary (*i.e.*, Lagrangian) action along a given trajectory is the instantaneous Lagrangian energy L(t).

As a final piece of notation which will help us generalize our results to the time-dependent case, we will sometimes write U'(t) to refer to the "instantaneous" unitary transformation that applies over an infinitesimal time interval dt at time t, that is,

$$U'(t) :\equiv U(t \to t + dt)$$

= 1 + iH(t)dt. (16)

Note also that any larger transformation $U(t_1 \rightarrow t_2)$ can be expressed as the timeordered product of all the infinitesimal U'(t) over the continuum of times t in the range from t_1 to t_2 . That is, we can write

$$U(t_1 \to t_2) = \mathcal{T} \prod_{t=t_1}^{t_2} U'(t)$$
 (17)

with the opposite ordering if $t_2 < t_1$. Thus, U'(t) uniquely defines U(t), so we will sometimes refer to U'(t) as the unitary trajectory also.

We should keep in mind that although the complete unitary trajectory U(t)(or U'(t)) between t_1 and t_2 determines the overall transformation $U(t_1 \rightarrow t_2)$, the converse is not true: Knowing the cumulative $U = U(t_1 \rightarrow t_2)$ for a particular pair of times t_1, t_2 is of course insufficient to determine a unique unitary trajectory U(t), since in general infinitely many cumulative action operators $A = A(t_1 \rightarrow t_2)$ can exponentiate to yield the same cumulative U (since expression (13) is multivalued), and furthermore, in the time-dependent case, a continuum of different Hamiltonian trajectories H(t) (which determine U'(t)) could implement a given cumulative action operator A.

We will similarly use the notation A'(t) = H(t)dt to denote the infinitesimal action operator that applies from time t to t + dt; note that $U'(t) = e^{iA'(t)} = 1 + iH(t)dt$.

6 Defining Computational Effort

With the above general definitions and observations aside, let us now proceed to define our concept of the amount of computational effort exerted by a system in undergoing a state trajectory $|\psi(t)\rangle$ between two times.

We will find it easiest to define this quantity first for the case of a system with a time-independent Hamiltonian H(t) = H = const. Later, we will show how our results can be generalized to the time-dependent case.

Let $|v\rangle$ be any eigenvector of H, and ω the corresponding eigenvalue, which is real since H is Hermitian. That is, let $H|v\rangle = \omega |v\rangle$. Thus, $|v\rangle$ is also an eigenvector of the cumulative action operator A(t) = Ht for any t, with eigenvalue $\alpha = \omega t$.

First, when t is an infinitesimal dt, consider the instantaneous U' = 1 + iHdt. Clearly, $|v\rangle$ is an eigenvector of U', since $U'|v\rangle = (1+iHdt)|v\rangle = (1+i\omega dt)|v\rangle = u|v\rangle$, where the scalar $u = 1 + i\omega dt = e^{i\omega dt} = e^{id\alpha}$. Thus, under application of U', the eigenvector $|v\rangle$ transforms to $|v'\rangle :\equiv e^{i\omega dt}|v\rangle = e^{id\alpha}|v\rangle$, that is, it phase-rotates in the complex plane at angular velocity ω through an infinitesimal angle $d\alpha$. Note also that

$$\Im \langle v | v' \rangle = \Im \langle v | (1 + \mathrm{id}\alpha) | v \rangle = \Im (1 + \mathrm{id}\alpha) \langle v | v \rangle$$
$$= \mathrm{d}\alpha = \langle v | \omega \mathrm{d}t | v \rangle = \langle v | A' | v \rangle = A'[v].$$
(18)

That is, when $|v\rangle$ is an eigenvector of H, the magnitude of the imaginary part of the inner product between infinitesimally adjacent state vectors is equal to the expectation value A'[v] of the infinitesimal action operator A' = H dt applied to the state. As we go on, we will extend the relationship (18) to non-infinitesimal trajectories, non-eigenvectors, and time-dependent Hamiltonians.

Next, note that the eigenvectors $|v\rangle$ of H are also eigenvectors of the cumulative action operators A(t) = Ht and cumulative unitaries $U(t) = e^{iA(t)} = e^{iHt}$, and vice-versa. Let $A(t)|v\rangle = \alpha(t)|v\rangle$, with $|v\rangle$ a fixed eigenket of A(t), and with $\alpha(t) = \omega t$ as its eigenvalue. Then, $U(t)|v\rangle = e^{iA(t)}|v\rangle = e^{i\alpha(t)}|v\rangle = u(t)|v\rangle$ where $u(t) = e^{i\alpha(t)}$. Thus, upon the application of U, $|v\rangle$ gets multiplied by the phase factor u(t), or (we can say) rotated by a total phase angle of $\alpha(t) = \omega t$, which could be much greater than 2π in long evolutions, as can also be seen by integrating $d\alpha$ over t. Note also that if we integrate $\Im \langle v | v' \rangle$ along the trajectory, we still get the cumulative action A(t)[v(0)]:

$$\int_{\tau=0}^{t} \Im\langle v(\tau) | v'(\tau) \rangle = \int_{\tau=0}^{t} \Im\langle v(\tau) | (1 + i\omega d\tau) | v(\tau) \rangle$$
(19)

$$= \omega t = \alpha(t) = \langle v(0) | A(t) | v(0) \rangle.$$
(20)

Next, consider an arbitrary pure state $|\psi(0)\rangle = \sum_i c_i(0)|v_i\rangle$, where the $|v_i\rangle$ are normalized eigenstates of H with eigenvalues ω_i , and the $c_i(0)$ are the initial coefficients of the $|v_i\rangle$ in the superposition. The state at time t can be expressed as

$$\begin{split} |\psi(t)\rangle &= \sum_{i} \exp[i\alpha_{i}]c_{i}(0)|v_{i}\rangle \\ &= \sum_{i} \exp[i\omega_{i}t]c_{i}(0)|v_{i}\rangle \\ &= \sum_{i} c_{i}(t)|v_{i}\rangle, \end{split}$$
(21)

where we see that each coefficient $c_i(t) = \exp[i\omega_i t]c_0(t)$ (in the fixed basis $\{|v_i\rangle\}$) simply phase-rotates with angular velocity ω_i along an origin-centered circle in the complex plane with constant radius $r_i = |c_i|$. Over any amount of time t, we see that c_i rotates in the complex plane by a total angle of $\alpha_i = \omega_i t$, while the line in the complex plane that joins c_i to the origin sweeps out an arc with an area of $a_i = \frac{1}{2}\omega_i tr_i^2$. (See fig. 3 for an illustration of the area swept out in the infinitesimal case.) For example, in time $t = 2\pi/\omega_i$, coefficient c_i sweeps out a complete disc of area $a_i = \pi r_i^2$ as it traverses an angle of $\alpha = 2\pi$. For consistency, in the case of clockwise rotations (negative ω_i), we will consider the area swept out to also be negative.

Now, let $\psi'(t) = \psi(t + dt)$. Then

$$\int_{\tau=0}^{t} \Im\langle\psi(\tau)|\psi'(\tau)\rangle = \int_{\tau=0}^{t} \Im\sum_{i} \bar{c}_{i}(\tau)c_{i}(\tau + \mathrm{d}\tau)$$
(22)

$$= \int \sum_{i} r_i^2 \Im\{ \mathrm{e}^{-\mathrm{i}\theta_i(\tau)} \mathrm{e}^{\mathrm{i}[\theta_i(\tau) + \omega_i \mathrm{d}\tau]} \}$$
(23)

$$= \int \sum_{i} p_i \Im(1 + \mathrm{i}\omega_i \mathrm{d}\tau) \tag{24}$$

$$= \int \sum_{i} p_i \mathrm{d}\alpha_i \tag{25}$$

$$= \int d\alpha = \alpha(t) = A(t)[\psi(0)]$$
 (26)

where the overbar denotes complex conjugation, $r_i = |c_i|$ as before, $\theta_i(\tau) = \arg(c_i(\tau))$, and α is now the weighted-average value of α_i .

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Now, consider the *total* area a(t) swept out by *all* coefficients c_i over time t. Note that $r_i^2 = |c_i|^2$ is also the probability p_i of basis state v_i , and so the *total*

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Fig. 3 In the energy eigenbasis, a complex coefficient c_i of a basis state sweeps out a small wedge-shaped area (shown exaggerated) in the complex plane over an infinitesimal time interval dt.

area swept out is always exactly half of the *average* angle $\alpha(t)$ of phase rotation (weighted the by state probability), or in other words, half of the expectation value of the A(t) operator applied to the state $\psi(0)$. That is,

$$a(t) = \sum_{i} \frac{1}{2} \omega_{i} t r_{i}^{2}$$

$$= \frac{1}{2} \sum_{i} p_{i} \alpha_{i}$$

$$= \frac{1}{2} A(t) [\psi(0)] = \frac{1}{2} \alpha(t).$$
 (27)

Thus we have shown that for time-independent Hamiltonians, the expectation value of the action operator A(t) applied to any initial state $\psi(0)$ is equal to the integral over the state trajectory of the inner product between infinitesimally adjacent states $\psi(t)$ and $\psi'(t) = \psi(t + dt)$ along the trajectory, as well as to the average phase angle α accumulated and to twice the complex-plane area *a* swept out by the state's coefficients, when the state is decomposed in the energy eigenbasis.

Of course, the inner product between two state vectors is a pure geometric quantity, and so is basis-independent. Therefore, the integral of $\Im\langle\psi|\psi'\rangle$ over the state trajectory does not depend at all on the (fixed) choice of basis under which states are decomposed into components. Likewise, the operator A(t) itself is a geometric object not inherently associated with any particular basis. Therefore,

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the identity

$$\int_{\tau=0}^{t} \Im\langle\psi(\tau)|\psi'(\tau)\rangle = A(t)[\psi(0)]$$
(28)

that we proved above is a fundamental one whose truth does not rely on any particular basis or coordinate system.

However, it is perhaps somewhat less obvious that the average angle α of phase rotation and the complex-plane area *a* swept out by the state coefficients should also be basis-independent quantities, since their original definitions explicitly invoked a choice of basis (the energy basis). However, in the next section we will show that in fact, these quantities are basis-independent as well. Thus, all of the following identities still hold true, regardless of basis:

$$2a = \alpha = \int_{\tau=0}^{t} \Im\langle\psi|\psi'\rangle = A(t)[\psi(0)], \tag{29}$$

where a is the total complex-plane area swept out by the state coefficients in any fixed basis, $\alpha = \int \omega dt$ is the time-integral of the expected value ω of the angular velocity ω_i of the state coefficients in any fixed basis, $\psi = \psi(\tau)$ is the state trajectory, with $\psi' = \psi(\tau + d\tau)$, A(t) is the action operator as we defined in equation (10), and we are using our mean-value notation $A(t)[\psi(0)] = \langle \psi(0) | A(t) | \psi(0) \rangle$.

Our proposed measure of the amount of change undergone (and computational effort exerted) along a state trajectory $\psi(t)$ generated by a constant H will then just be the α value for that trajectory.

Later, in section 8, we will show that the above identities also still hold even when H(t) varies over time, and so our measure will generalize to that case as well.

7 Generalizing to Arbitrary Bases

The above discussion made use of a set of basis vectors $\{|v_i\rangle\}$ which were taken to be orthonormal eigenvectors of the (temporarily presumed constant) Hamiltonian operator H. Now, we will show that this particular choice of basis was in fact unnecessary, and that the same statements concerning the relationship between the area swept out, the average phase angle accumulated, and the action A(t) would remain true in any fixed (time-independent) basis.

At first, it may seem very non-obvious that the area swept out should still be exactly half of the action. Note that our previous arguments for this relied on the fact that in the energy basis $\{|v_i\rangle\}$, the coefficients c_i all rotate at uniform angular velocities ω_i in circles in the complex plane, while their individual magnitudes remain constant. In a different basis $|v_j\rangle$ (distinguished by using a different index symbol j), this will no longer be true. Each basis vector $|v_j\rangle$ in the new basis is in general some superposition of the $\{|v_i\rangle\}$, such as

$$|v_j\rangle = \sum_i u_j^i |v_i\rangle,\tag{30}$$

where the matrix $\mathbf{U} = [u_j^i]$ of complex coefficients (with the subscript j indexing rows, and the superscript i indexing columns) is, most generally, any unitary matrix. We can also write this equation in matrix-vector form as $|v_j\rangle = \mathbf{U}|v_i\rangle$, where the over-arrow here denotes that we are referring to the entire column-ordered sequence of basis vectors, $|v_i\rangle = \begin{bmatrix} |v_1\rangle \\ \vdots \end{bmatrix}$. Of course, a general state vector ψ can equally well be expressed as a linear superposition of either set of basis vectors, that is,

$$|\psi\rangle = \sum_{i} c_{i} |v_{i}\rangle \tag{31}$$

$$|\psi\rangle = \sum_{j} c_{j} |v_{j}\rangle. \tag{32}$$

But now, we can substitute eq. (30) into eq. (32) and rearrange, as follows:

$$|\psi\rangle = \sum_{ij} c_j u_j^i |v_i\rangle = \sum_i \left(\sum_j c_j u_j^i\right) |v_i\rangle.$$
(33)

Now, since the $|v_i\rangle$ are linearly independent, the expansion of $|\psi\rangle$ in terms of them must be unique, so we can equate the coefficients on $|v_i\rangle$ in equations (31) and (33) to get

$$c_{i} = \sum_{i} u_{j}^{i} c_{j}$$
$$\overrightarrow{c_{i}} = \mathbf{U}^{\mathrm{T}} \overrightarrow{c_{j}}, \qquad (34)$$

where T is matrix transpose. We can easily solve this equation for the c_j coefficients as follows:

$$\overline{c_i} = \mathbf{U}^{\mathrm{T}} \overline{c_j}$$

$$(\mathbf{U}^{\mathrm{T}})^{-1} \overline{c_i} = \overline{c_j}$$

$$\overline{\mathbf{U}} \overline{c_i} = \overline{c_j}$$

$$c_j = \sum_i \overline{u_j}^i c_i.$$
(35)

In other words, each complex coefficient in the new basis is just a particular linear combination of what the various complex coefficients were in the old basis.

If the coefficients c_i in the old energy basis are describing perfect circles around the complex origin at a variety of radii and angular velocities, there is no guarantee that the coefficients c_j in the new basis will still be describing circular paths centered on the origin, although their paths will of course still be continuous and smooth if the original c_i trajectories were. In general, the c_j will follow complicated looping trajectories in the complex plane, generated as if by Ptolemaic planetary epicycles, *i.e.*, as a sum of circularly rotating vectors. A given c_j will in general return to its initial location in the complex plane only when its components Fig. 4 Area swept out (exaggerated) by a coefficient c_j (in a basis other than the energy eigenbasis) over an infinitesimal time interval dt. Note that both its phase and its magnitude change, in general.

 c_i that have nonzero values of u_j^i all simultaneously return to their initial locations exactly, which might even take infinitely long, if the corresponding ω_i values were relatively irrational.

Anyhow, the important point for our present purposes is that the c_j s do not, in general, maintain a constant magnitude (distance from the origin), and so the area swept out by the c_j over a given time is no longer just a section of a circle, which was very easy to analyze. Instead, while c_j 's phase angle θ_j is rotating, simultaneously its magnitude r_j may also be growing or shrinking. Fig. 4 illustrates the situation.

To clarify what we mean by the phase angle $\theta_j(t)$ a bit more carefully, let us use $d\alpha_j(t) \approx 0$ to denote the infinitesimal increment of phase angle from times t to t + dt such that

$$d\alpha_j \equiv \arg(c'_j) - \arg(c_j) \pmod{2\pi},\tag{36}$$

so that $d\alpha_j$ remains infinitesimal even when c_j crosses a branch cut of the Arg() function. Then, let $\alpha_j(t)$ be the total accumulated phase angle over time t, that is, the integral of $d\alpha_j$ over time,

$$\alpha_j(t) = \int_{\tau=0}^t \mathrm{d}\alpha_j \tag{37}$$

so that $\alpha_j(0) = 0$. Now, just let $\theta_j(t) = \operatorname{Arg}[c_j(0)] + \alpha_j(t)$. Thus also $d\theta_j = d\alpha_j$.

What, now, is the area swept out in our new basis? First, notice that in the infinitesimal limit, it is exactly half of the area of the parallelogram that is spanned on two adjacent sides by $c_j = c_j(t)$ and $c'_j = c_j(t + dt)$, considered as vectors in the complex plane. See fig. 5.

The parallelogram area, itself, is $da_j = r_j r'_j \sin(d\theta_j)$, where r_j and r'_j are the magnitudes of the old and new coefficients, respectively. However, note that the

Fig. 5 The infinitesimal area da_j swept out approaches one-half of the parallelogram area $r_j r'_j \sin d\theta_j$.

area da_j of this parallelogram is also the signed magnitude of the scalar "cross product" $c_j \times c'_j$ between the coefficients, considered as vectors in the complex plane. (The traditional cross product, defined in three dimensions, would be a vector perpendicular to the complex plane having this value da_j as its length.) There is a nice identity [20] connecting the scalar cross product and dot product with the conjugate multiplication of complex numbers, namely:

$$\bar{c}d = c \cdot d + i(c \times d), \tag{38}$$

where \bar{c} means the complex conjugate of c, and $c \cdot d$ denotes the real scalar "dot product" between c and d considered as vectors, namely $|c||d| \cos[\arg(d) - \arg(c)]$, and $c \times d$ denotes the real scalar "cross product" previously mentioned, namely $|c||d| \sin[\arg(d) - \arg(c)]$.

Applying this identity to our situation, we can see that the area swept out, since it is half the cross product, is half of the imaginary part of the conjugate product $\bar{c}_j c'_i$ between the old and new coefficients, and also to half of $\sin(d\alpha_j) = d\alpha_j$;

$$\mathrm{d}a_j = \frac{1}{2}\mathrm{d}\alpha_j = \frac{1}{2}\Im(\bar{c}_j c'_j). \tag{39}$$

Now, this is just the area swept out by a single component c_j . To find the total area da swept out by all coefficients, we merely sum over components:

$$da = \frac{1}{2} \sum_{j} \Im(\bar{c}_{j}c'_{j}) = \frac{1}{2} \Im \sum_{j} \bar{c}_{j}c'_{j}$$
$$= \frac{1}{2} \Im\langle\psi|\psi'\rangle = \frac{1}{2} d\alpha$$
(40)

In other words, *just like in the energy basis*, in an arbitrary basis, it is still true that the infinitesimal increment da in the area swept out by the coefficients is exactly one-half of $\Im\langle\psi|\psi'\rangle$, the imaginary component of the inner product between infinitesimally adjacent vectors $\psi = \psi(t)$ and $\psi' = \psi(t + dt)$ along the trajectory, and further that this is equal to half of $d\alpha = d\theta$, the average increment of the continuously-varying phase angles $\theta_j(t)$ of the coefficients.

Now, we saw earlier that $\Im\langle\psi|\psi'\rangle$ is also equal to the expectation value $A'[\psi] = \langle\psi|A'|\psi\rangle$ of the infinitesimal action operator A' = Hdt applied to the state ψ , for any state ψ . So in connection with the result (40) that we just obtained, this means

that $A'[\psi]$ gives *exactly* the average phase angle accumulation $d\alpha$ of the coefficients c_j of ψ in *any* basis, and twice the complex-plane area da swept out by those coefficients. We can thus think of A' as being the operator representation of a fundamental, basis-independent concept of "average angle accumulated" or "total area swept out" over infinitesimal intervals.

8 Generalizing to Time-dependent Hamiltonians

In the previous section, we established the basis-independence of the identities $2da = d\alpha = \Im\langle \psi | \psi' \rangle = \omega dt = A'[\psi] = \langle \psi | H dt | \psi \rangle$ for infinitesimal changes of the state vector ($\psi \rightarrow \psi'$) along its trajectory over infinitesimal time intervals dt, under any *constant* Hamiltonian H.

But, as long as the Hamiltonian H(t) only changes in continuous fashion, it can always be considered essentially "constant" throughout any infinitesimal interval dt, even if it is varying over non-infinitesimal timescales. Therefore, the above identities will still hold true instantaneously even for a time-dependent Hamiltonian H(t), which is what we originally started out our discussion with. Thus, when we integrate the above equation (40) over time, it remains true that:

$$2a = \alpha = \int_{t=t_1}^{t_2} \Im\langle\psi(t)|\psi(t+\mathrm{d}t)\rangle \tag{41}$$

$$= \int_{t=t_1}^{t_2} \omega(t) \mathrm{d}t \tag{42}$$

$$= \int_{t=t_1}^{t_2} \langle \psi(t) | H(t) | \psi(t) \rangle \mathrm{d}t \tag{43}$$

$$= \int_{t=t_1}^{t_2} A'(t)[\psi(t)].$$
(44)

In words, this says that for any initial state ψ , we have that 2a (twice the complexplane area swept out by the coefficients of ψ , in any basis) is equal to α , the average phase angle swept out by the state coefficients, as well as to (41) the integral along the trajectory $\psi(t)$ of the imaginary component of the dot product between neighboring vectors along the trajectory, and also to (42) the integral of the average phase velocity of the coefficients, weighted by the instantaneous basis state probabilities $p_i(t) = r_i(t)^2$, which is (43) the time-integral of the instantaneous Hamiltonian energy $E(t) = H(t)[\psi(t)]$ of the instantaneous state $\psi(t)$, which (finally) is (44) the integral of the infinitesimal actions $d\alpha(t) = \langle \psi(t) | A'(t) | \psi(t) \rangle$ on the instantaneous states $\psi(t)$.

The natural next question to ask is, given that $A'[\psi] = d\alpha$ remains true over infinitesimal intervals dt in the general time-dependent case, and given that cumulatively, $A(t)[\psi(0)] = \alpha$ in the time-independent case (H(t) = H = const.), does this cumulative relation still hold true in the general time-dependent case? That is, for A(t) (as defined in eq. (10)) is it still true that

$$A(t)[\psi(0)] = \alpha \tag{45}$$

even if the phase angle α was accumulated under the influence of a varying H(t)?

If this equation (45) is universally correct, then we will have a very nice, simple interpretation for the general action operator A(t) even in the case of a timedependent H(t), namely that, when applied to any initial state $\psi(0)$, it simply gives the angular length α of the trajectory that will be traversed by that state, a quantity which obeys all of the identities (41)-(44).

Actually it seems that this is true, and the proof is quite elegant. First, from eq. (17) and the boundary condition U(0) = 1, fix U = U(t), the overall unitary transform operating between times 0 and t that is implied by the values of the time-dependent Hamiltonian $H(\tau)$ for all $0 \le \tau \le t$. Fix then also A = A(t) by using eq. (13) and the associated discussion, using the continuity requirement on $A(\tau)$ and the requirement that A(0) = 0.

Now, consider any eigenvector $|\phi_i\rangle$ of U, which is a state that undergoes a cyclic evolution (in the projective Hilbert space) under $H(\tau)$ or any other process (Hamiltonian trajectory) that implements U, since $U|\phi\rangle = \mu_i |\phi_i\rangle$, with μ_i being the associated unit-modulus eigenvalue. Of course, $|\phi_i\rangle$ is then also an eigenvector of A, with an eigenvalue α_i such that $A|\phi_i\rangle = \alpha_i |\phi_i\rangle$ and $\mu_i = e^{i\alpha_i}$.

To see that this α_i must indeed be the same as the total phase angle α accumulated by $|\phi_i\rangle$ as defined in *e.g.* eq. (44), consider that once the overall operator A has been determined, we can simply divide it by t to find an alternative *time-independent* $H_c = A/t$ that would also generate the very same action operator A and the same unitary U when applied over the same time interval t. From the discussion in section 6, is is easy to see that the value of α is then indeed exactly the phase angle accumulated from the initial state $|\phi_i\rangle$ when implementing A via this (alternative) time-independent H_c .

Now, does *every* Hamiltonian trajectory that implements A (including our original time-dependent $H(\tau)$) involve the same total accumulation α of phase angle? We can see that it must, because any trajectory $H(\tau)$ can, it seems, be continuously deformed into the constant trajectory $H_c(\tau) = H_c$ while maintaining the same overall A (and thus U) throughout the deformation process. At no point during this continuous deformation process can the total phase α that is accumulated ever change, since, to produce the same U, the total phase α must always remain congruent to $\alpha_i \pmod{2\pi}$, and it would be impossible for the total phase accumulated to jump by a multiple of 2π at any point during any continuous deformation of the trajectory.

To see that this is true, recall from eq. (13) and the associated discussion that any continuous $A(\tau)$ can be characterized by a continuously varying eigenbasis $\{|u_i(\tau)\rangle\}$ of $U(\tau)$ (with a sort of k-dimensional continuous gauge freedom, where k is the Hilbert space dimension), and by implied integer parameters $n_i(\tau)$ that select which of the logarithm values must be used at each time point τ . As we continuously deform the Hamiltonian trajectory $H(\tau)$ as well as the eigenbases $\{|u_i(\tau)\rangle\}$ (and thus the gauges of the associated eigenvalues $u_i(\tau)$), the set of time points τ at which the $n_i(\tau)$ values change also changes continuously. Nowhere during this continuous, local process can the total angle α accumulated along the trajectory possibly change discontinuously by a multiple of 2π . On the Interpretation of Energy as the Rate of Quantum Computation

Thus, our arbitrary time-dependent $H(\tau)$ takes the eigenstate $|\phi_i\rangle$ through the same total angle α as would the constant H_c for which we already know that $\langle \phi_i | A | \phi_i \rangle = \alpha$.

The above discussion establishes that (regardless of the dynamics H(t)) the A operator that we derive from it always gives the correct accumulated angle α for all eigenstates ϕ_i of A; therefore it is also correct for arbitrary initial superposition states $\psi(0)$ (and for mixed states as well).

For a final interesting observation, let $\alpha(\psi(0), t)$ denote the angle α accumulated from the initial state $|\psi(0)\rangle$ over time t, and note that since

$$\langle \psi(0)|A(t)|\psi(0)\rangle = \alpha(\psi(0), t) \tag{46}$$

for all initial $\psi(0)$, the time-derivative of the operator A(t) must satisfy

$$\langle \psi(0) | \frac{\mathrm{d}}{\mathrm{d}t} A(t) | \psi(0) \rangle = \frac{\partial}{\partial t} \alpha(\psi(0), t).$$
(47)

Recall meanwhile that $d\alpha(t)$ is given by applying A'(t) = H(t)dt to the state $\psi(t)$; *i.e.*, $d\alpha(t) = A'(t)[\psi(t)]$. Of course, $\psi(t) = U(t)\psi(0)$, so we have that

$$\langle \psi(0) | \frac{\mathrm{d}A}{\mathrm{d}t}(t) | \psi(0) \rangle = \frac{A'(t)}{\mathrm{d}t} [U(t)\psi(0)] \tag{48}$$

$$= \langle \psi(0) | U^{\dagger}(t) H(t) U(t) | \psi(0) \rangle.$$
(49)

and thus

$$\frac{\mathrm{d}A}{\mathrm{d}t}(t) = U^{\dagger}(t)H(t)U(t)$$
$$= \mathrm{e}^{-\mathrm{i}A(t)}H(t)\mathrm{e}^{\mathrm{i}A(t)}.$$
(50)

Now, note that applying the time-dependent operator form (6) of the Schrödinger equation to $U(t) = e^{iA(t)}$, we get

$$\frac{\mathrm{d}}{\mathrm{d}t}\mathrm{e}^{\mathrm{i}A(t)} = \mathrm{i}H(t)\mathrm{e}^{\mathrm{i}A(t)}$$

$$= \mathrm{i}\mathrm{e}^{\mathrm{i}A(t)}\mathrm{e}^{-\mathrm{i}A(t)}H(t)\mathrm{e}^{\mathrm{i}A(t)}$$

$$= \mathrm{e}^{\mathrm{i}A(t)}\frac{\mathrm{d}}{\mathrm{d}t}[\mathrm{i}A(t)],$$
(51)

where we have used (50) in the last step. In other words, the ordinary rule $de^f = e^f df$ for the differential of an exponential of a function f actually turns out to be true when f = iA(t), despite the fact that the Hamiltonian may be time-dependent and that A(t) doesn't necessarily even commute with its time-derivative! This is due to the special way in which we defined our A(t) function, and would not be true for more general time-dependent operators.

9 Discussion of Effort

Although a choice of a particular cumulative action operator A still gives us freedom to choose any number of different Hamiltonian trajectories $H(\tau)$ for implementing it, over various total amounts of time t, we have seen above that all such trajectories are equivalent in terms of the total amount α of phase angle that is accumulated starting from any fixed initial state $|\psi(0)\rangle$.

As hinted previously, we might even consider the quantity α (or, more properly, its absolute value) to be a reasonable definition of the geometric *length* of the path that a normalized state vector $|\psi(t)\rangle$ describes as it moves along any continuous path (parameterized by any real variable t) along the unit sphere in Hilbert space, since (note) α depends only on the shape of the state trajectory itself, and not on any other properties of the Hamiltonian trajectory, such as the energy of other orthogonal states.

As a result, an intrinsic metric on the normalized Hilbert space is provided by the distance function

$$d(|\psi_1\rangle, |\psi_2\rangle) :\equiv \min |\alpha| \tag{52}$$

where α is the accumulated phase angle along a given trajectory, and the minimum is taken over all normalized, continuous paths from $|\psi_1\rangle$ to $|\psi_2\rangle$, or a subset of such that is deemed available. The absolute-value operator is required in order to obtain a proper (positive) metric, since trajectories with unboundedly negative values of α could exist if we allow states to have negative energy. Paths having the minimum absolute α between a given pair of states can be considered to be (sections of) geodesics on the normalized Hilbert space.

In [21], Wootters introduced a statistically-motivated distance metric between quantum states which he called "statistical distance," and showed that it was identical to the ordinary Hilbert-space distance function $d(\psi_1, \psi_2) = \arccos |\langle \psi_1 | \psi_2 \rangle|$. It turns out that our distance function d above is in fact exactly the same as this also, if all Hilbert-space trajectories are considered. However, if the space of allowed trajectories is restricted (for example, if the Hamiltonians are forced to be local) then a different distance measure results. In Wootters' metric, the distance between any two distinguishable states (*e.g.*, two different randomly chosen computational basis states) is only $\arccos 0 = \pi/2$, whereas if we define distance by minimizing over allowed trajectories, we could obtain a much greater figure.

Later, we will see that our distance measure will also allow us to derive a natural metric on unitary operations, telling us the "distance" between two unitaries, as measured by the difficulty of getting from one to the other, in terms of the minimum distance traversed by worst-case states.

Anyway, noting that this measure α of trajectory length which we have explored above is stable with respect to changes of basis, that there are multiple simple ways of defining it, and that it connects strongly with fundamental physical concepts such as action and energy, as well as with primitive geometric concepts such as angles and areas, and that it forms a natural metric on the Hilbert space, all of these facts together motivate us to propose this measure as being the most natural and genuine measure of the total "amount of change" that is undergone by

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a physical quantum state vector $|\psi(t)\rangle$ as it changes dynamically under a (possibly varying) physical influence H(t).

Insofar as we can consider *all* dynamical evolution and change to be forms of "computation," where this word is construed in a very general sense, we can also accept this measure as being an appropriate measure of the *amount of computational effort exerted* by the system as it undergoes the given trajectory.

Thus, from here on, rather than calling our quantity "action" (which would lead to confusion with the action of the Lagrangian), or "accumulated phase angle" (which is awkward) we will refer to our quantity as simply the *effort* when we wish to be concise, and abbreviate it with the symbol \mathcal{F} . That is,

$$\mathcal{F}_{t_1 \to t_2} \left[\psi(t) \right] :\equiv \int_{t=t_1}^{t_2} \Im\langle \psi(t) | \psi'(t) \rangle \tag{53}$$

is a real-valued functional of a state vector trajectory $\psi(t)$ taken between two times t_1 and t_2 . Note that the value of \mathcal{F} depends only on the shape of the path. It is independent of the absolute time, the speed at which the trajectory is traversed, and on various other details of the Hamiltonian that generates the trajectory (such as its eigenvalues for eigenstates that are not components of ψ); in general, many different Hamiltonian evolutions can generate the same path, which will always have the same total effort. So, in the above equation, we can consider $\psi(t)$ to just be a parameterized curve where t is now just any *arbitrary* real-valued parameter, not necessarily even corresponding to physical time. In other words, the *effort* quantity does not depend on the precise system of coordinates that is used for measuring the passage of time, but rather only on a pure geometric object, namely the path taken through Hilbert space.

Note that to say that the path length corresponds to computational *effort* is not to imply that all of the physical computation that is occurring in the given system is necessarily being harnessed and applied by humans to meet our calculational needs, only that this is the total amount of raw computational work that is occurring "in nature." The choice of the word "effort" is intended to evoke the commonsense realization that effort may be wasted, *i.e.*, not used for anything useful.

Note also that the action operator A (as we have defined it) gives a concise yet particularly comprehensive characterization of a given computational process, in the sense that it determines not only the overall unitary operation $U = e^{iA}$ that will be performed, but also the amount of effort that will be expended in getting to the final result from any given initial state.

The primary caveat to the above conception of computational effort seems to be that the quantity \mathcal{F} (together with the rate of phase rotation, and the path length in Hilbert space) is dependent on where we choose to draw our zero of energy. As is well known, absolute energies are only physically defined up to an additive constant, and so the total Hamiltonian action or effort is only well defined up to this constant multiplied by the elapsed time t.

A natural and widely-used convention is to define the least eigenvalue of the Hamiltonian (the "ground state" energy) to be the zero of energy. In a similar fashion, we can choose to additively shift the Hamiltonian so that the least eigenvalue

of the cumulative action operator A(t) is taken to represent zero effort. (Note that this approach can even be used when the Hamiltonian itself is time-dependent.)

However, this choice is by no means mandated mathematically, and in fact, in certain pathological cases (such as an infinite-dimensional or time-dependent Hamiltonian with unboundedly negative eigenvalues), there might not even be *any* minimum eigenvalue for the resulting action operator over a given interval. One needs to keep these caveats in the back of one's mind, although they seemingly end up not very much affecting the potential practical applications of this concept, which we will address in a later section.

Another reason that we might not want to consider the ground state energy to always be zero is if the ground state energy varies, especially if it includes energy that had to be explicitly transferred into the system from some other external subsystem. Thus, energy that is present in a given system, even if that system is in its ground state, may still represent energy that was transferred from elsewhere and isn't being used for other purposes; *i.e.*, it may represent "wasted" computational effort, and we may wish to count it as such, rather than just counting it as zero effort.

Another possible convention would be to count a system's energy as being its total (gravitating) mass-energy, or rest mass-energy, if we want it to be independent of the observer's velocity. One might think this choice is a somewhat less arbitrary than the ground state convention, since mass is a physical observable, but unfortunately, in general relativity, the contribution to the total mass-energy of a local system that is due to its gravitational self-energy isn't actually independent of the coordinate system that is used ([22], p. 62). However, this caveat is usually only important in extreme systems such as neutron stars and black holes, where the gravitational self-energy contributes significantly to the system's total mass.

In any case, for now, we propose to just make a "gentlepersons' agreement" that we will always make sure that the energy eigenvalues of the systems that we consider are always shifted so as to be positive, so that the total effort is always positive, and we don't have to worry about what would be the meaning of a negative "amount of computational effort." Unfortunately, this strategy rules out considering certain classes of systems, such as bottomless potential wells, or the infinite Dirac sea of negative-energy fermion states. But resolving this issue will have to wait for future work.

10 More Abstract Scenarios

In the above, we have specified a well-defined (at least, up to an additive constant) positive, real-valued measure \mathcal{F} of the amount of computational effort represented by any trajectory of a state vector in Hilbert space.

This raises the question of whether we can assign a measure of computational effort to other physical situations that may be less completely specified. For example, we may be given a cumulative action operator A, but not know the detailed Hamiltonian trajectory $H(t)|_{t=t_1}^{t_2}$ that generated it, and we may be given only a set V of possible initial states (rather than a single definite state), or we may have

a probability distribution or density function $p: V \rightarrow [0, 1]$ over initial states. In such more abstract situations, can we still meaningfully define the amount of computational effort exerted by the system as it undergoes the evolution specified by its Hamiltonian over a given time interval?

Of course we can. Given a cumulative action operator A and given any specific state $\psi = \psi(t_1)$ at the initial time t_1 , the value of $\mathcal{F}_{t_1 \to t_2}[\psi(t)]$ is independent of the details of the Hamiltonian trajectory H(t) and is given simply by

$$\mathcal{F}_A(\psi) :\equiv A[\psi] = \langle \psi | A | \psi \rangle, \tag{54}$$

which can be called *the effort undergone by* ψ *under* A.

We can therefore also naturally express the *average or expected effort over* V *exerted by the action operator* A as:

$$\widehat{\mathcal{F}}_{V}(A) = \operatorname{Ex}_{V}[\mathcal{F}_{A}] = \sum_{\psi \in V} p(\psi)\mathcal{F}_{A}(\psi) = \langle A \rangle = \operatorname{Tr}(\rho A),$$
(55)

where the density operator ρ describing the initial mixed state is constructed from the probability distribution over pure states ψ in the usual fashion, that is, with $\rho = \sum_{\psi \in V} p(\psi) |\psi\rangle \langle \psi|$. If no probability distribution p has been provided, we can use a uniform distribution over some natural measure on the set V.

This then gives us a workable definition of the mean effort exerted by a system over time under a given Hamiltonian, even when the initial state is not exactly known.

In some situations, we might also be particularly interested in the *maximum* effort over the set V of possible initial states. For example, suppose we are preparing the initial state of the system, and we want to initialize the system in such a way that it will exert the maximum effort possible. Given A and maximizing over V, we define the *maximum effort exerted by* A over V as

$$\mathcal{F}_{V}^{+}(A) :\equiv \max_{\psi \in V} \mathcal{F}_{A}(\psi).$$
(56)

This can be considered to be a measure of the *potential* computational "strength" of the given action operator A, expressing that any Hamiltonian H(t) that implements A over some arbitrary interval $t_1 \rightarrow t_2$ could exert an amount $\mathcal{F}_V^+(A)$ of computational effort over that same interval, given a suitable initial state. Insofar as the actual state that we end up getting *might* be the one that undergoes the maximal amount of effort, we can say that a system with an unknown or unspecified state is, at least, exerting this much "potential" computational effort.

Even if the actual state turns out *not* to be the maximal-action one, the system could still be thought of as having "done the work" of determining that the actual state is *not* the one that should have transitioned through the given maximum Hilbert-space distance. This particular thought should really be credited to Seth Lloyd, who pointed out to me in personal discussions, as an analogy, that an ordinary Boolean gate operation can still be thought of as doing computational work even if the output bit that it is applied to is not actually changed; namely, it is doing the work of determining *that* the bit should not change.

Similarly to how we defined the maximum effort, we can likewise define the *minimum effort of A over V* as $\mathcal{F}_V^-(A) :\equiv \min_{\psi \in V} \mathcal{F}_A(\psi)$, although we should keep in mind that if the ground state of the action operator A is an available initial state in V, and if we use the convention that the ground state action is defined to be zero, then $\mathcal{F}_V^-(A)$ will always be 0, and so will not be very useful.

11 Difficulty of Performing an Operation

Suppose now that we are given *no* information about the situation to be analyzed except for a unitary operator U on the Hilbert space \mathcal{H} , and we want to address the following question: How much computational effort, at minimum, is required to physically implement U? By "implement" we mean that U is the time evolution operator that ends up being generated by the dynamics over some interval, according to $U = e^{iA}$ for some action operator A. We can call this minimum required effort the *difficulty* \mathcal{D} of implementing the unitary operator U. Our framework gives us a natural way to formalize this notion.

Assuming we have some freedom of choice in the design of the system, then among the set \mathcal{A} of all Hermitian operators A on \mathcal{H} , or among at least a set $\aleph \subseteq \mathcal{A}$ of *available* or implementable action operators, we might want to choose the operator A that generates U that has the *smallest* value of the maximum or worstcase effort $\mathcal{F}_V^+(A)$ over the set V of possible initial state vectors. This A can be considered to be the "best" action operator for generating the given unitary U, in the sense that the length of the longest trajectory that would be undergone by any possible state vector $\psi \in V$ is minimized. This strategy is analogous to what we do in traditional algorithm design, where we usually choose the algorithm that has the minimum time complexity on worst-case input data. In our case, A can be considered to abstractly represent the algorithm selected, while the initial vector ψ represents the input data. Rather than time complexity, we focus on effort or Hamiltonian action, since (as we will see) this translates directly to time when a given supply of energy is available to be invested in the system.

In some situations, it may be preferred to choose A so as to minimize the *expected* effort rather than the worst-case effort, for example, if we want to minimize the total effort exerted over an arbitrarily large set of computations with randomly chosen input states selected from some distribution.

We can thus define the maximum $(\mathcal{D}^+_{\aleph,V})$ and expected $(\widehat{\mathcal{D}}_{\aleph,V})$ difficulty of a desired unitary transform U under the available action set \aleph and initial-state set V as follows:

$$\mathcal{D}^{+}_{\aleph,V}(U) :\equiv \min_{A \in \aleph} \mathcal{F}^{+}_{V}(A)$$
$$= \min_{A \in \aleph} \max_{\psi \in V} \mathcal{F}_{A}(\psi)$$
(57)

$$\widehat{\mathcal{D}}_{\aleph,V}(U) :\equiv \min_{A \in \aleph} \widehat{\mathcal{F}}_V(A)$$
$$= \min_{A \in \aleph} \sum_{\psi \in V} p(\psi) \mathcal{F}_A(\psi)$$
(58)

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Note that in all cases we still want to minimize over the available action operators $A \in \aleph$, because there is usually no physical reason why indefinitely large action operators (which waste arbitrarily large amounts of effort) could not be constructed to implement a given unitary; thus, maximizing over action operators would thus always give ∞ and would not be meaningful.

A remark about the set ℵ of available action operators. Typically it would be constrained by what constitutes an "available" dynamics that we are free to choose within a given theoretical, experimental, or manufacturing context. For example, ℵ might reasonably be constrained to include only those action operators that are obtainable from time-dependent Hamiltonians H(t) which are themselves constructed by summing over local interaction terms between neighboring subsystems, or by integrating a Hamiltonian density function that includes only local terms on a field over some topological space, e.g., to reflect the local structure of spacetime in a quantum field theory picture. Or, we might constrain ourselves to action operators that are obtainable from time-independent Hamiltonians only, e.g. if we are designing a self-contained (closed) quantum system. Finally, practical considerations may severely constrain the space of Hamiltonians to ones that can be readily constructed in devices that can be built using a specific manufacturing process, although we should note that if scalable universal quantum computers can be built, then any desired local Hamiltonian could be straightforwardly emulated on these machines.

As a brief aside, it is also interesting to note that a given difficulty function $\mathcal{D}(U)$ (either the worst-case or average-case version, and whatever \aleph and V are) also induces an intrinsic metric on the space of unitaries of a given rank; we can define a suitable distance function between unitaries by

$$d(U_1, U_2) = \mathcal{D}(U_2 U_1^{\dagger}) \tag{59}$$

that is, the distance between U_1 and U_2 in this metric is just the difficulty of performing the relative unitary $U_{1\rightarrow 2} :\equiv U_2 U_1^{\dagger}$ that is equivalent to undoing U_1 (using $U_1^{\dagger} = U_1^{-1}$) and then doing U_2 . A unitary trajectory for implementing $U_{1\rightarrow 2}$ that actually minimizes the effort will then form, when right-multiplied by U_1 , a (section of a) geodesic in the space of unitaries passing between the unitaries U_1 and U_2 (since $U_{1\rightarrow 2}U_1 = U_2$). Of course, in general, the shortest unitary trajectory for implementing $U_{1\rightarrow 2}$ will *not* actually work by doing U_1^{\dagger} followed by U_2 ; for example, if U_1 and U_2 have high difficulty but are very close together, then the shortest unitary trajectory between them will be much more direct than this.

Now, given our notion of the computational difficulty of a given unitary U, we can now reinterpret previous results (such as [5,16]) regarding "quantum speed limits" or minimum times to implement various specific unitary transforms of interest, or classes of transforms, given states of specified average energy above the ground state, as follows: These analyses are implicitly specifying an \aleph (usually, just all Hermitian operators) and a V (usually, just the entire Hilbert space), and showing that the worst-case difficulty $\mathcal{D}^+(U)$ for the transform U has a specific value (or lower bound), assuming the presence of a time-independent Hamiltonian where the ground state energy is usually set to 0. In other words, such analyses

show that a certain minimum worst-case effort or Hamiltonian action is required to implement the particular U in question.

As an example, Margolus and Levitin's result [5] can be interpreted as telling us that any U that rotates some state ψ to an orthogonal state has a worst-case difficulty of $\mathcal{D}^+(U) \ge h/4$, since their result shows that any state of energy E takes time at least h/4E (no matter what the Hamiltonian) to accumulate the action needed to take it to an orthogonal state; thus the Hamiltonian action A = Et that is required to carry out such a transition is at least h/4.

Another result in [5] implies that if there is a ψ such that $(|\psi\rangle, U|\psi\rangle, U^2|\psi\rangle$, $\dots, U^{N-1}|\psi\rangle, U^N|\psi\rangle = |\psi\rangle$) comprises a cycle of N states, with each orthogonal to the preceding and succeeding states in the cycle, then $\mathcal{D}^+(U) \geq \frac{h}{2} \frac{N-1}{N}$, even if we are given complete freedom in constructing the Hamiltonian, aside from a requirement that it be time-independent. For N = 2, this expression reduces to h/4, while for $N \to \infty$, it goes to h/2. Thus, any physical computation that proceeds autonomously though an unbounded sequence of distinct states must exert at least h/2 effort per state transition.

Notice that the Margolus-Levitin theorem is, strictly speaking, only giving us a *lower bound* on the worst-case difficulty, since it is considering only a particular state ψ of interest (namely, one that actually undergoes a transition to an orthogonal state), rather than finding the worst-case potential effort to perform the corresponding U, maximized over all possible initial ψ in the Hilbert space. Later, we will see that the actual worst-case effort for an orthogonalizing transformation is actually $h/2 = \pi$ even in the N = 2 case, and possibly even higher in cases that go through more states.

We anticipate that, armed our definitions, it would be a highly useful and worthwhile exercise to systematically go through a variety of the quantum unitary transforms that have already been identified in quantum computing as comprising useful "quantum logic gate" operations, and quantify their worst-case and average difficulty, according to the above definitions, under various physically realistic sets of constraints. This would directly tell us how much physical Hamiltonian action is required to carry out those operations (given a best-case Hamiltonian implementation, while operating on a worst-case or average-case input state). We can likewise do the same for classical reversible Boolean logic operations embedded within unitary operations, as well as classical irreversible Boolean logic operations embedded within classical reversible operations, with ancilla bits used as needed for carrying away garbage information to be discarded.

Such an investigation will, for the first time, give us a natural and physically well-founded measure of the physical complexity of logic operations, in terms of Hamiltonian action. This in turn would directly tell us the minimum physical time to perform these operations within any physical system or subsystem using a set of states having a given maximum energy about the ground state, given the known or prespecified constraints on the system's initial state and its available Hamiltonian dynamics. This new quantification of computational complexity may also allow us to derive lower bounds on the number of quantum gates of a given type that would be required to implement a given larger transformation in terms of smaller ones,

and possibly to show that certain constructions of larger gates out of smaller ones are optimal.

In subsequent subsections, we begin carrying out the above-described line of research, with some initial investigations of the difficulty of various simple operations in situations where the available dynamics is relatively unconstrained, which is the easiest case to analyze.

12 Specific Operations

In this section, we explore the difficulty (according to our previous definitions) of a variety of important quantum and classical logic operations.

We will begin by considering some educated guesses about the difficulty of various unitaries. For each unitary U we are to imagine implementing it via a particular transformation trajectory U'(t) (and Hamiltonian H(t) such that $U'(t) = e^{iH(t)dt}$) that is as "direct" as possible, in the sense of minimizing the Hilbert-space distance through which worst-case states are transported. Intuition tells us that these minimal trajectories are expected to follow geodesics in the space of unitaries, as per the metric we defined earlier; in other words, they should be "straight-line" paths, so to speak, that get us to the desired unitary as directly as possible.

12.1 General two-dimensional unitaries

Let us begin by considering U_2 , the space of unitary transformations on Hilbert spaces of dimensionality 2. In quantum computing, these correspond to singlequbit quantum logic gates. As is well known (*e.g.*, see [23], eq. 4.9), any such U can be decomposed as

$$U = e^{i\alpha} R_{\hat{n}}(\theta) \tag{60}$$

where $\hat{n} = (n_x, n_y, n_z)$ is a real 3D unit vector and $R_{\hat{n}}(\theta)$ is a Bloch-sphere rotation about this vector by an angle of θ , that is,

$$R_{\hat{n}}(\theta) = e^{i(\theta/2)(\hat{n}\cdot\boldsymbol{\sigma})} \tag{61}$$

where $\boldsymbol{\sigma} = (\sigma_x, \sigma_y, \sigma_z)$ is the vector of Pauli matrices

$$\sigma_x = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}, \sigma_y = \begin{bmatrix} 0 & -\mathbf{i} \\ \mathbf{i} & 0 \end{bmatrix}, \sigma_z = \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix}.$$
 (62)

Let us now consider breaking down U into its multiplicative factors $e^{i\alpha}$ and $R_{\hat{n}}(\theta)$, which we observe commute with each other, since $e^{i\alpha}$ is a scalar. Thus, we can consider these two components of U to be carried out in either order, or even simultaneously if we prefer.

Let's start by looking at $R_{\hat{n}}(\theta)$. At first, we might guess that the worst-case effort that is required to perform $R_{\hat{n}}(\theta)$ for angles θ where $-\pi \leq \theta \leq \pi$ ought to just turn out to be $|\theta|/2$, since, for example, a Bloch sphere rotation through an angle of $\theta = \pi$ radians corresponds to inverting a spin in ordinary 3D space

through an angle of 180° to point in the opposite direction, which is an orthogonalizing transformation, and we already know from the Margolus-Levitin theorem that any transition to an orthogonal state under a constant Hamiltonian requires a minimum action (given zero ground state energy) for the state in question of $h/4 = (\pi/2)\hbar = (\pi/2)$ rad, or an area swept out of $\pi/4$ square units. This is a good first guess, but later, we will see that the actual worst-case action turns out to be twice as large as this. (Our intuition forgot to take into account the fact that the state vector in the Margolus-Levitin theorem isn't actually the worst-case one, as far as the accumulated Hamiltonian action is concerned.)

Indeed, for any real unit 3-vector \hat{n} (the "axis of rotation" for the Bloch sphere), one can easily verify that there is always a corresponding complex state vector

$$|v_{\hat{n}}^{+}\rangle = \frac{1}{\sqrt{2(1+n_{z})}} \begin{bmatrix} n_{z}+1\\ n_{x}+\mathrm{i}n_{y} \end{bmatrix}$$
 (63)

which is a unit eigenvector of $\hat{n} \cdot \boldsymbol{\sigma}$ having eigenvalue +1. This state vector is therefore also an eigenstate of $R_{\hat{n}}(\theta)$, with eigenvalue $e^{i(\theta/2)}$. In other words, in any orthonormal basis that includes $|v_{\hat{n}}^+\rangle$ as one of the basis vectors, as θ increases from 0 (for now, we'll assume for simplicity that the final value of θ is non-negative, $0 \leq \theta \leq \pi$), the coefficient of the $|v_{\hat{n}}^+\rangle$ component of the state $|\psi(t)\rangle = R_{\hat{n}}(\theta)|v_{\hat{n}}^+\rangle$ (starting from the initial state $|\psi(0)\rangle = |v_{\hat{n}}^+\rangle$, where the coefficient $c_{|v_{\hat{n}}^+\rangle}$ is 1) describes a circular arc in the complex plane centered on the origin, sweeping out a total angle of $\theta/2$, and an origin-centered area of $\theta/4$. As we saw earler, this same measure of the weighted-average accumulated angle and total area accumulated still holds in any basis. So, we have that the effort of $R_{\hat{n}}(\theta)$ must be at least $\theta/2$. Indeed, this is the exact worst-case effort, since $|v_{\hat{n}}^+\rangle$'s eigenvalue is maximal, so no pure energy eigenstate can possibly sweep out a larger angle as θ increases, and therefore no superposition of energy eigenstates (*i.e.*, no general state) can do so either.

Now, what about the $e^{i\alpha}$ factor that's included in the expression for a general $U \in U_2$? Note that this term represents an overall (global) phase factor that applies to all eigenstates. As such, even the ground state $|g\rangle$ of whatever Hamiltonian is used to implement U might still accumulate a phase due to this phase factor. In this case, $|g\rangle$ would have nonzero Hamiltonian energy. If we redefine $|g\rangle$ to instead have zero energy $(H|g\rangle = 0)$, then $|g\rangle$'s coefficient would not phase-rotate at all, since the action operator A = Ht would give $A|g\rangle = 0$ for this state, and $U|g\rangle$ would give $(e^{iA})|g\rangle = (e^0)|g\rangle = |g\rangle$, that is, $|g\rangle$ would be unchanged by this U. However, it does not follow that we can always just let α be zero, as $|g\rangle$ may generally have accumulated an additional phase resulting from the $R_{\hat{n}}(\theta)$ component of U as well. It is the *total* phase accumulated by the ground state that we wish to define to be zero.

Let us now consider the following: Under the transformation $R_{\hat{n}}(\theta)$, as θ increases from 0, we notice that $|v_{\hat{n}}^+\rangle$ (the eigenvalue-1 eigenstate of $\hat{n} \cdot \boldsymbol{\sigma}$ which we constructed above) only phase-rotates by an angle $\theta/2$. Under $U = e^{i\alpha}R_{\hat{n}}(\theta)$, $|v_{\hat{n}}^+\rangle$ therefore undergoes an overall phase-rotation by an angle of $\alpha + \theta/2$. We confidently conjecture that the "least potential action" or most efficient way to implement U is to apply a Hamiltonian that simultaneously sweeps both α and θ

forward steadily from 0, at respective rates that are exactly proportional to their intended final values. If this is correct, then $|v_{\hat{n}}^+\rangle$ is indeed an eigenstate of that best-case Hamiltonian, with energy $(\alpha + \theta/2)/t$ (recall that we're using $\hbar = 1$), where t is the total time taken for α and θ to reach their final values.

However, since the space we are working with is two-dimensional, there must be another energy eigenstate as well. Solving the eigen-equation $(\hat{n} \cdot \boldsymbol{\sigma})|v\rangle = r|v\rangle$, we find that the other eigenvalue r of $\hat{n} \cdot \boldsymbol{\sigma}$ is -1, and the other unit-length eigenvector, modulo phase-rotations, is (for $n_z > 0$)

$$|v_{\hat{n}}^{-}\rangle = \frac{1}{\sqrt{2(1-n_z)}} \begin{bmatrix} n_z - 1\\ n_x + in_y \end{bmatrix}$$
 (64)

or, in the special case when $n_z = 0$, then instead any normalized column vector $|v_{\hat{n}}^-\rangle = [v_0; v_1]$ where $|v_0| = |v_1| = 2^{-1/2}$ will work, so long as the vector components v_0 and v_1 have the specific obtuse (that is, $> 90^\circ$) relative phase angle that is given by the relation $v_1 = (-n_x - in_y)v_0$. (Note that $|n_x + in_y| = 1$ when $n_z = 0$.)

Thus, for any Hamiltonian that smoothly sweeps θ forward in a steady transformation $R_{\hat{n}}(\theta)$ with $\theta \propto t$, there will actually be two different energy eigenstates having energies that are negatives of each other, one state in which the accumulated action of the Hamiltonian is $\theta/2$ (as we saw above), and another state (the ground state) where the action is the negative of this, or $-\theta/2$. Together with the global phase-rotation of α , we have that the total action for U is $\alpha + \theta/2$ and $\alpha - \theta/2$ for these two energy eigenstates, respectively.

Following our convention that the total action in the ground state should be always considered to be zero, we can shift the energy levels upwards in such a way that the lower value $\alpha - \theta/2$ will be equal to 0, in other words, we can adjust our rate of global phase rotation (which determined α) in such a way that we have exactly $\alpha = \theta/2$. Now, the total action in the high energy state is $\alpha + \theta/2 = \theta/2 + \theta/2 = \theta$.

In other words, starting with any $U \in U_2$ and decomposing it as $U = e^{i\alpha}R_{\hat{n}}(\theta)$, which involves a rotation of the Bloch sphere through an angle of θ about an axis \hat{n} , we can calculate a meaningful difficulty $\mathcal{D}^+(U)$ by using the convention that the ground state should be considered to have energy 0, and by letting $\mathcal{D}^+(U) = \mathcal{D}^+(U_{\hat{n}}(\theta))$, where we define $U_{\hat{n}}(\theta) \equiv e^{i\theta/2}R_{\hat{n}}(\theta)$, that is, ignoring the original value of α (whatever it was) and instead adjusting α to have the value $\alpha = \theta/2$ which assigns the ground state to zero energy. Thus, we can say that the "true" computational/physical difficulty of U (given this choice) is exactly θ for any single-qubit unitary $U = e^{i\alpha}R_{\hat{n}}(\theta)$, regardless of the value of α . If θ is a pure number (implicitly bearing an angle unit of radians), then the worst-case Hamiltonian action to carry out the desired transform using the best-case Hamiltonian (assuming that is indeed what we have managed to characterize above) is $\theta\hbar$, in whatever physical units we wish to express \hbar . That is, $\mathcal{D}^+(U) = \theta$.

To wrap up this section, let us take a look at the precise form of the Hamiltonian that we are proposing. Note that

$$\hat{n} \cdot \boldsymbol{\sigma} = \begin{bmatrix} n_z & n_x - \mathrm{i}n_y \\ n_x + \mathrm{i}n_y & -n_z \end{bmatrix}$$
(65)

is itself an Hermitian operator which plays the role of the Hamiltonian operator H with respect to the Bloch-sphere rotation unitary $R_{\hat{n}}(\theta) = e^{i(\theta/2)(\hat{n}\cdot\sigma)}$, if the rotation angle θ is taken be equal to twice the time t. Meanwhile, in this scenario, the extra phase-rotation factor $e^{i\alpha} = e^{i(\theta/2)}$ out front corresponds simply to an additional constant energy of +1, using the same angular velocity units of $(\theta/2t)$. This gives us a total "Hamiltonian" (in quotes because we haven't introduced an explicit time parameter here yet) of $H_{\hat{n}}$ that is required to implement a steady rotation about \hat{n} which is equal to

$$H_{\hat{n}} = 1 + \hat{n} \cdot \boldsymbol{\sigma}$$

$$= \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} + \begin{bmatrix} n_{z} & n_{x} - in_{y} \\ n_{x} + in_{y} & -n_{z} \end{bmatrix}$$

$$= \begin{bmatrix} 1 + n_{z} & n_{x} - in_{y} \\ n_{x} + in_{y} & 1 - n_{z} \end{bmatrix}.$$
(66)

With this choice of "Hamiltonian," we can easily check that the $|v_{\hat{n}}^{\pm}\rangle$ are indeed its energy eigenstates, with $H_{\hat{n}}|v_{\hat{n}}^{-}\rangle = 0$ (the ground state has "energy" 0) and $H_{\hat{n}}|v_{\hat{n}}^{+}\rangle = 2$, which is what we want since it will cancel out with the 2 in the denominator of the exponent in the rotation unitary $U_{\hat{n}}(\theta) = e^{i\theta/2}R_{\hat{n}}(\theta) = e^{i(\theta/2)(1+\hat{n}\cdot\sigma)} = e^{i(\theta/2)H_{\hat{n}}}$.

To generalize the picture slightly, if a rotation through θ about an axis \hat{n} is to take place over an arbitrary amount of time t, then we require a Hamiltonian (a proper one now, in actual angular-velocity energy units) of

$$H = \frac{\theta}{2t}H_{\hat{n}} = \frac{\theta}{2t} \begin{bmatrix} 1 + n_z & n_x - in_y \\ n_x + in_y & 1 - n_z \end{bmatrix}$$
(67)

With this choice of Hamiltonian, note that things works out nicely so that the high-energy eigenstate $|v_{\hat{n}}^+\rangle$ phase-rotates at exactly the desired rate $\omega^+ = \theta/t$, since we have that

$$H|v_{\hat{n}}^{+}\rangle = \frac{\theta}{2t}H_{\hat{n}}|v_{\hat{n}}^{+}\rangle = \frac{\theta}{2t}2|v_{\hat{n}}^{+}\rangle = \frac{\theta}{t}|v_{\hat{n}}^{+}\rangle = \omega^{+}|v_{\hat{n}}^{+}\rangle.$$
 (68)

Thus, the action operator A = Ht comes out exactly equal to the angle operator Ω which gives the total angle of phase rotation for both the energy eigenstates $|v_{\hat{n}}^{\pm}\rangle$, that is, $A|v_{\hat{n}}^{-}\rangle = \Omega|v_{\hat{n}}^{-}\rangle = 0|v_{\hat{n}}^{-}\rangle$ and $A|v_{\hat{n}}^{+}\rangle = \Omega|v_{\hat{n}}^{+}\rangle = \theta|v_{\hat{n}}^{+}\rangle$. And for an arbitrary initial state ψ , *i.e.*, for any normalized complex superposition of the eigenstates $|v_{\hat{n}}^{\pm}\rangle$, $A[\psi] = \Omega[\psi]$ gives the quantum mean angle of phase rotation.

Note that in all the above discussion, we have assumed that the rotation angle is non-negative, *i.e.*, that $0 \le \theta \le \pi$ (rad). To complete the picture, note that for values of θ between 0 and $-\pi$, we can convert them to positive angles by the simple expedient of rotating instead by an angle of $|\theta| = -\theta$ about the $-\hat{n}$ axis, which is an exactly equivalent rotation. This has the effect of exchanging the values of the $|v_{\hat{n}}^{\pm}\rangle$ eigenstates, as well as the sign of the $H_{\hat{n}}$ component of H. Other than that, everything else is the same, with the result that the action A always comes out non-negative and equal to the absolute value of θ . Of course, for the case of absolute

angles outside the range $(-\pi, \pi]$, we can just reduce them to the equivalent angle in $(-\pi, \pi]$ by adding or subtracting the appropriate multiple of 2π .

In the above, although we have not yet quite finished proving rigorously that the specific H we have given is in fact the one that implements U with the least possible value of the worst-case action A, still, we expect that it should already seem highly plausible to the reader that this should in fact be the case, due to the directness and simplicity of our construction, which made use only of the simple fact that any arbitrary $U \in U_2$ can be decomposed into a single generalized rotation about an arbitrary axis is real three-space, accompanied by a global phase rotation. Of course, a more complete proof of the optimality of this construction would be desirable to have, but it will have to wait for future work.

12.2 Specific single-qubit gates

Given the above discussion, to determine the difficulty \mathcal{D} of any single-qubit gate U is a simple matter of finding some unit 3-vector \hat{n} and angles $\alpha, \theta \in (-\pi, \pi]$ such that $U = e^{i\alpha}R_{\hat{n}}(\theta)$, which is always possible. This then establishes that $\mathcal{D}^+(U) = |\theta|$, under our ground zero energy convention. Let us look briefly at how this calculation comes out for various single-qubit gates of interest.

- 1. The Pauli spin-operator "gates" $X = \sigma_x$ (which is the in-place NOT operation in the computational basis), $Y = \sigma_y$, and $Z = \sigma_z$ all of course involve a rotation angle of $\theta = \pi$, since they all square to the identity $(2\pi \text{ rotation})$. Thus, $\mathcal{D}^+(X) = \mathcal{D}^+(Y) = \mathcal{D}^+(Z) = \pi = h/2$.
- 2. The "square root of NOT" gate $N = \frac{1}{2} \begin{bmatrix} 1+i & 1-i \\ 1-i & 1+i \end{bmatrix}$ of course requires an angle of $\pi/2$, since $N^2 = X$. Thus, $\mathcal{D}^+(N) = \pi/2 = h/4$.
- 3. The Hadamard gate $N = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & -1 \\ 1 & -1 \end{bmatrix}$ requires a rotation angle of π about the $\hat{n} = (1, 0, 1)/\sqrt{2}$ axis, *i.e.*, $\hat{n} \cdot \boldsymbol{\sigma} = (\sigma_x + \sigma_z)/\sqrt{2}$. Also note that $H^2 = 1$ and a rotation through 2π is the identity. Thus, $\mathcal{D}^+(H) = \pi = h/2$.
- 4. The "phase gate" $S = \begin{bmatrix} 1 & 0 \\ 0 & i \end{bmatrix}$ requires $\theta = \pi/2$ since note that $S^2 = Z$. So, $\mathcal{D}^+(S) = \pi/2 = h/4$.
- 5. The so-called " $\pi/8$ " gate $T = \begin{bmatrix} 1 & 0 \\ 0 & \exp[i\pi/4] \end{bmatrix}$ involves $\theta = \pi/4$ since note that $T^4 = Z$. Thus, $\mathcal{D}^+(T) = \pi/4 = h/8$.
- 6. The generalized phase gate $ph(\theta) = \begin{bmatrix} 1 & 0 \\ 0 & \exp[i\theta] \end{bmatrix}$ is just a rotation by an angle of θ about the *z* axis, so $\mathcal{D}^+(ph(\theta)) = \theta = \theta\hbar$.

As a point of comparison, the paper [16] studies the time required to perform the specific gate $U = e^{i\theta}X$ (*i.e.*, NOT with global phase rotation) using an optimal Hamiltonian, and conclude that the minimum time τ required (for a specific initial state) is

$$\tau = \frac{h}{4E} \left(1 + 2\frac{\theta}{\pi} \right). \tag{69}$$

Note that the corresponding Hamiltonian action α or effort \mathcal{F} is

$$\alpha = \mathcal{F} = E\tau = \frac{h}{4} + 2\frac{h}{4}\frac{\theta}{\pi}$$
$$= \frac{\pi}{2}\hbar + \theta\hbar$$
$$= \frac{\pi}{2} + \theta \quad (\text{with } \hbar = 1).$$
(70)

At first glance, this might appear to contradict our claim that the difficulty of such a U ought to be exactly π . However, we should keep two things in mind. First, in [16], Levitin *et al.* are concerned with the time to carry out U in the case of a specific subset of initial states which will actually transition to an orthogonal state in the time τ . However, these particular states are not the "worst-case" ones from our perspective, and so they don't determine the maximum effort. Rather, the particular states under consideration in their paper all have a mean energy of only $\overline{E} = (E_1 + E_2)/2$, where E_1 and E_2 are the low and high energy eigenvalues of the ideal Hamiltonian, respectively. Letting $E_1 = 0$ (our ground zero assumption), we have that $E_2 = 2E$. Since E_2 has the highest energy available given this spectrum, the E_2 energy eigenstate accumulates more action over the time τ than any other possible state, in particular, double that of states with energy $\overline{E} = E_2/2$, and thus it is the E_2 state that determines the worst-case action, which is twice that of [16], or in other words $A = \pi$. The term involving θ in (70) drops out entirely, since as we already saw earlier, global phase shifts are irrelevant when considering total action, under our convention that the ground state action is always defined to be zero. Levitin *et al.* don't make this adjustment, because they are assuming that the Hamiltonian has already been arranged in advance to have a desired energy scale. Thus, the global phase rotation by θ leads to an extra additive θ in their expression (70) for the action.

12.3 Difficulty of achieving infidelity

A natural and widely-used measure of the degree of closeness or similarity between two quantum states u, v is the *fidelity*, which is defined (for pure states) as $F(u, v) = |\langle u | v \rangle| = |u^{\dagger}v|$. (See [23].) Note that if the actual state of a system is u, and we measure it in a measurement basis that includes v as a basis vector, the square of the fidelity $p = F^2$ gives the probability that the measurement operator will project the state down to v, and that v will be seen as the "actual" state. (This is a "quantum jump" or "wavefunction collapse" event, or, in the many-worlds picture, it is the subjectively experienced outcome when the state of the observer becomes inextricably entangled with that of the system.) Likewise with the roles of u and v reversed. Thus, only when F = 0 are the states u and v orthogonal.

We can also define a related quantity, the "infidelity" $Inf(u, v) \equiv \sqrt{1-p} = \sqrt{1-F^2}$. The squared infidelity between u and v is then just the probability 1-p that if the actual state is u, then it will *not* be taken to v by a projective measurement (in a measurement basis that includes v), and vice-versa. In other words, if v is some old state of a system, and u is its new state, the squared infidelity between

u and v is the probability that the answer to the question "Is the state different from v yet?" will be found to be "yes" when this question is asked experimentally by a measurement apparatus that compares the state with v.

Let us now explore the minimum effort that is required in order for some of the possible state vectors of a system to attain a given degree of infidelity (relative to their initial states), in the case of two-dimensional Hilbert spaces. Note that not all vectors will achieve infidelity; in particular, the eigenvectors of any timeindependent Hamiltonian will always have 0 infidelity.

We start by recalling from earlier that any 2-dimensional unitary can be considered a rotation of the Bloch sphere about some axis in ordinary (real-valued) 3-D space. Since a simple change of basis suffices to transform any axis to any other, we can without loss of generality presume a rotation about the z axis, represented by

$$R_{\hat{z}}(\theta) = \begin{bmatrix} e^{-i\theta/2} & 0\\ 0 & e^{i\theta/2} \end{bmatrix}.$$
 (71)

We saw earlier that the effort of any such rotation (under the ground-zero convention) is always exactly θ . What initial state will gain infidelity most rapidly under this transformation? Until we figure this out, let us allow the initial state to be a general unit vector $|v\rangle = [v_0; v_1] = v_0 |0\rangle + v_1 |1\rangle$ in the basis $|0\rangle, |1\rangle$. Then $|u\rangle = R_{\hat{z}}(\theta)|v\rangle = [e^{-i\theta/2}v_0; e^{i\theta/2}v_1]$ as a column vector of complex coefficients. Now the fidelity between v and u is

$$F(v, u) = |\langle v|u \rangle| = |\langle v|R_{\hat{z}}(\theta)|v \rangle| = |v_0^* e^{-i\theta/2} v_0 + v_1^* e^{i\theta/2} v_1| = |e^{-i\theta/2}|v_0|^2 + e^{i\theta/2}|v_1|^2| = \left| \left[\cos\frac{\theta}{2} - i\sin\frac{\theta}{2} \right] |v_0|^2 + \left[\cos\frac{\theta}{2} + i\sin\frac{\theta}{2} \right] |v_1|^2| = \left| \left(\cos\frac{\theta}{2} \right) (|v_0|^2 + |v_1|^2) + i\left(\sin\frac{\theta}{2} \right) (|v_1|^2 - |v_0|^2) \right| = \left| \left(\cos\frac{\theta}{2} \right) + i\left(\sin\frac{\theta}{2} \right) (|v_1|^2 - |v_0|^2) \right|.$$
(72)

where in the last line we have made use of the fact that $|v_0|^2 + |v_1|^2 = 1$ for a normalized v. Now, F^2 is the sum of the squared real and imaginary components of the expression inside the outermost absolute-value delimiters || above:

$$[F(u,v)]^{2} = \Im^{2}[\langle v|u\rangle] + \Re^{2}[\langle v|u\rangle]$$

= $\cos^{2}\left(\frac{\theta}{2}\right) + \sin^{2}\left(\frac{\theta}{2}\right)\left(|v_{1}|^{2} - |v_{0}|^{2}\right)^{2}$
= $\cos^{2}\left(\frac{\theta}{2}\right) + \sin^{2}\left(\frac{\theta}{2}\right)\left(1 - 4|v_{1}|^{2}|v_{0}|^{2}\right)$
= $1 - 4\sin^{2}\left(\frac{\theta}{2}\right)|v_{1}|^{2}|v_{0}|^{2},$ (73)

where in getting from the second to the third line, we have again made use of the fact that $|v_0|^2 + |v_1|^2 = 1$. We can reassure ourselves that the last line of (73) is always in the range [0,1], since $|v_0|^2|v_1|^2 \le 1/4$ given that $|v_0|^2 + |v_1|^2 = 1$. Note also that the fidelity is minimized when $|v_0|^2 = |v_1|^2 = \frac{1}{2}$, that is, when the two *z*-basis states are in an equal superposition. This is then the "worst case" (worst in terms of "least fidelity") which we wish to focus on.

So now, the infidelity $I = Inf(u, v) = \sqrt{1 - F^2(u, v)}$ comes out to be a reasonably simple expression:

$$Inf(u,v) = \sqrt{1 - [F(u,v)]^2} = \sqrt{4\sin^2\left(\frac{\theta}{2}\right)|v_1|^2|v_0|^2}$$
(74)

$$= 2\left(\sin\frac{\theta}{2}\right)|v_0||v_1|. \tag{75}$$

Note that for any given angle of rotation in $0 < \theta < \pi/2$, the infidelity is maximized when $|v_0| = |v_1| = 1/\sqrt{2}$. For such v, we have $|v_0||v_1| = \frac{1}{2}$ and so

$$Inf(u,v) = \sin\frac{\theta}{2}.$$
(76)

Thus, if we wish that some system initially in state v should achieve a desired degree I of infidelity (relative to its initial state) using a transformation of minimum effort, we must choose a unitary transformation that is a rotation $R_{\hat{n}}(\theta)$ about an axis \hat{n} that is "perpendicular" to v, and rotate by an angle $\theta = 2 \cdot \arcsin(I)$. The Hamiltonian action α accumulated by "worst-case" (that is, maximum-energy) vectors under this transformation is (by definition) the difficulty $\mathcal{D}^+(R_{\hat{n}}(\theta))$ of that unitary, and is given by $\alpha = 2 \cdot \arcsin(I)$.

However, the specific initial vector v that we are dealing with will not have the maximum energy E (relative to ground) but rather half of this, or E/2, since half of its probability mass will be in the high-energy state, and half in the zeroenergy ground state. Therefore, v's total Hamiltonian action (amount of change) along its trajectory will instead be exactly $\alpha(v) = \arcsin(I)$, a wonderfully simple expression. This α is the effort exerted by the specific state v as it traverses a maximally efficient path for achieving infidelity $I = \sin \alpha$.

So, for example, suppose we want to cause some given initial state v to transition to a new state that has only a probability of at most p = 1/2 of being confused with the initial state if it were measured. This is to say that the infidelity between the states should be at least $I = \sqrt{1-p} = 1/\sqrt{2}$, which requires the state to traverse a trajectory that has a length of at least $\theta = \arcsin(I) = \arcsin(1/\sqrt{2}) = \pi/4 = h/8$, which can be done using a minimum-difficulty unitary transform whose worst-case effort is twice as great as this, or $\pi/2 = h/4$, meaning that the worst-case (maximum-energy) states of the system would traverse a trajectory of this (greater) length under an optimal implementation of such a transformation.

Assuming that the actual given initial state in question is assigned an average energy of only E above the ground state, it will take time at least t = h/8E to

carry out a unitary transformation on this state that achieves a probability above 1/2 of distinguishing it from the resulting state; whereas, if we are given that the *maximum* energy state in the qubit spectrum has energy E, then it will take time at least t = h/4E to carry out the transform.

In other words, to carry out an operation in time t that yields a 50% probability (or less) of conflation of some initial states with their successors requires that the initial states in question must have energy at least E = h/8t, and that states of energy at least E = h/4t must exist in the spectrum.

Note that the above results are also perfectly consistent with the Margolus-Levitin theorem [5]. That is, plugging in an infidelity of I = 1 to represent a transition to an orthogonal state, we find that the specific initial state's effort $\mathcal{F}(v) = \arcsin(1) = \pi/2$ while the worst-case difficulty for this transform is $\theta = 2 \arcsin(1) = \pi$; these figures are twice that for the previous example. And so for a state to attain a 0% probability of conflation (*i.e.*, to reach an orthogonal state) requires that it have at least twice the energy as the previous scenario, or $E = \pi/2t = h/4t$ (under the Hamiltonian used to carry out the transformation), while other energy levels of at least $\pi/t = h/2t$ must be present in the spectrum of the Hamiltonian operator being used.

12.4 Higher-dimensional operations

Naturally, we are interested not only in unitaries in U_2 , but also in higher dimensions, in particular, unitaries in the groups U_{2^n} , which correspond to general "quantum logic gate" operations (really, arbitrary quantum computations) operating on sets of n qubits.

In particular, let us focus on the "controlled-U" gates with one target bit, which take the general form (modulo qubit reorderings)

$$U' = \mathbf{C}^{n-1}U \equiv \begin{bmatrix} 1 & & \\ & 1 & \\ & \ddots & \\ & & U \end{bmatrix}$$
(77)

where we have $2^n - 2$ ones along the diagonal, and a rank-2 unitary matrix U in the lower-right corner. In other words, for computational basis states $|b_0b_1 \dots b_{n-1}\rangle$, whenever the first n - 1 qubits $b_0b_1 \dots b_{n-2}$ are not all 1's, the state remains unchanged; otherwise, the unitary U is performed on the final qubit b_{n-1} .

We observe immediately that $\mathcal{D}^+(U') \geq \mathcal{D}^+(U)$, since all the input states that undergo any change at all will undergo the exact same transformation (in the subspace associated with the last qubit) that they would if U were just applied unconditionally. Thus, the worst-case trajectories when conditionally applying Ucan be no shorter than the worst-case unconditional trajectories (under an optimal implementation).

Furthermore, if U by itself would be optimally implemented by the Hamiltonian H, then it is easy to believe that U' would likewise be optimally imple-

mented by the Hamiltonian

$$H' = \begin{bmatrix} 0 & & \\ 0 & & \\ & \ddots & \\ & & H \end{bmatrix}$$
(78)

that is, with 0's everywhere except for a copy of H in the lower-right 2×2 submatrix. It is easy to verify that this H', when exponentiated, indeed produces the desired U'. And since its worst-case difficulty is equal to our lower bound $\mathcal{D}^+(U)$, it is in fact an optimal H', assuming our earlier conjecture about the optimality of H is correct. In this case, if H' is actually an available Hamiltonian in the context one is considering, then the effort of U' is indeed exactly the same as the effort of U.

We can see from this example that when we consider the full space of mathematically describable Hamiltonians, we are likely to greatly underestimate the effort, compared to what can actually be implemented. The typical known implementations of U in terms of small local quantum gates would require a number of orthogonalizing operations that is at least linear in n, whereas in our case above, the effort is constant (upper-bounded by π). It seems likely that the effort for a physically realistic (*e.g.* field-theory based) Hamiltonian for this class of Us would have to be more than constant, since the interaction of n qubits to determine an outcome would appear to necessarily be a non-local process.

In most physical situations of interest, we will not necessarily have available Hamiltonians that are of any form desired, such as the form H' suggested above. Instead, we may only have available a more limited, perhaps parameterized suite of Hamiltonians, perhaps ones that are formed by a sum or time-sequence of specific, controllable, localized couplings having (say) at most 2 qubits each, as is popularly represented in the quantum computing literature using the schematic notation of quantum logic networks.

Obviously, whenever our space of available Hamiltonians is more restricted than the simple "all Hermitian operations" scenario analyzed above, the resulting values of $\mathcal{D}^+(U)$ will in general become much larger, and probably also much more difficult for us to analytically calculate. To compute $\mathcal{D}^+(U)$ for Hamiltonians that can plausibly be constructed within the context of particular experimental frameworks that are readily physically realizable in the lab (or in a manufactured product, *e.g.*, a someday-hopefully-to-be-realized commercial quantum computer) is clearly a much more complex and difficult task than we have attempted to tackle in this paper. To address this problem more fully will have to wait for future work.

Still, we hope that the present work can at least serve as a fruitful conceptual foundation on which we can proceed to build meaningful analytical and/or numerical analyses of the physical/computational "difficulty" of performing various quantum operations. We also hope that this work will serve as a helpful stepping stone for future investigators who wish to continue exploring the many deep and rich interconnections between physical and computational concepts.

12.5 Classical reversible and irreversible Boolean operations

Although in the above discussion we have focused on the effort required to carry out quantum gate operations, it is easy to extend the results to classical logic operations as well. Any classical reversible operation is just a special case of a quantum gate where the matrix elements of the unitary operator (in the computational basis) are 0 or 1. For example, a reversible Toffoli gate or Controlled-Controlled-NOT (CCNOT) is a special case of the C^2U gate addressed in §12.4 above. Specifically, since the U in question is X (NOT), which has a rotation angle of π , the effort required for Toffoli must be at least π , and indeed is exactly π if arbitrary Hamiltonians can be constructed. Toffoli is a universal gate for classical reversible computation, so a construction of any classical reversible circuit out of Toffoli gates sets an upper bound (as a multiple of π) on the difficulty of that computation, apart from any extra effort that may be required to control transitions between gates (which could be substantial, but is probably close to linear in the number of operations performed).

As for ordinary irreversible Boolean operations, these can be embedded into reversible operations as follows. Consider, for example, a standard boolean inverter, whose function is irreversible as it is normally specified in an electrical engineering context. The explicit function of an inverter is to destructively overwrite its output node with the logical complement of its input. (Please note that this function is distinct from that of a classical reversible NOT operation, which simply toggles a bit in-place.) Due to Landauer's principle, the physical information contained in the output node cannot actually be destroyed, but is instead transferred to reside in the environment. So, we can model the ordinary inverter's function as a sequence of reversible operations as follows:

- 1. Exchange output bit with an empty bit in the device's environment
- 2. Increment an "environment pointer" to refer to the next empty bit in some unbounded list
- 3. Perform a CNOT between input node and (now empty) output node

The first step can be understood as the emission from the device of the old stored value of the bit, in the form of entropy. The second step can be viewed as implementing the continuous flow of entropy away from the device, to make room for discarding the results of subsequent inverter operations. Finally, the third step carries out the desired logical function. The above breakdown is not necessarily the simplest possible implementation of the classical inverter (although it is probably close), but it at least sets an upper limit on the number of quantum operations that are absolutely required.

The first step can be carried out by a unitary SWAP operation between the two bits in question. The second step can be carried out by an annihilate/create pair of operations that moves a "particle" by one position to point to the next empty location in the environment; this corresponds to a unitary operation that increments the state vector $|i\rangle$ of some subsystem that specifies the integer location *i* of the environment pointer. Finally, the third step is just an ordinary CNOT, with an effort of π . In principle, we could calculate and add up the effort for all these steps, together with the effort needed to update a part of the machine state that keeps track of which step we are on, to arrive at an upper bound on the effort required to implement a classical inverter operation. However, this calculation might not be very meaningful unless we did more work to specify a detailed physical setup that would allow us to confirm that such a bound was achievable in a practical hardware implementation.

13 Relation to Berry phase

An interesting question to ask about our quantity \mathcal{F} is what relationship (if any) it has to the classic notion of the geometric or Berry phase of a quantum trajectory [24–31]. So far, the relationships between these concepts are not completely clear, and working them out in more detail will have to wait for future work. However, some initial remarks are in order.

Let H(t) be any time-dependent Hamiltonian that implements the unitary U for t going from 0 to τ , and let $|\psi\rangle$ be an eigenvector of U, with eigenvalue $e^{i\phi}$. The state $|\psi\rangle$ thus undergoes a cyclic evolution in the projective (phase-free) Hilbert space. Aharonov and Anandan [26] point out the relation $-\phi = \alpha - \beta$ (the integrated form of their equation (2)), where α is the integral of the instantaneous Hamiltonian energy of the state,

$$\alpha = \frac{1}{\hbar} \int_{t=0}^{\tau} \langle \psi(t) | H(t) | \psi(t) \rangle \mathrm{d}t$$
(79)

and β is a term given by

$$\beta = \int_{t=0}^{\tau} \langle \tilde{\psi}(t) | \mathbf{i} \frac{\mathrm{d}}{\mathrm{d}t} | \tilde{\psi}(t) \rangle \mathrm{d}t, \tag{80}$$

where $\tilde{\psi}(t)$ is any continuously gauge-twiddled version of $\psi(t)$ such that $\tilde{\psi}(0) = \tilde{\psi}(\tau) = \psi(0)$. Aharonov and Anandan's paper [26] revolves around their claim that this β quantity is a generalized version of the Berry phase that applies even to non-adiabatic evolutions.

However, if the results of the present paper are correct, then Aharonov and Anandan's β is always an arbitrary value congruent to 0 (modulo 2π) and thus is not a physically meaningful quantity. The reason is that the α in (79) is exactly our $\alpha = A[\psi(0)]$, where $U = e^{-iA}$ (in the usual sign convention, which A&A are using), and thus $\psi(0)$ is also an eigenvector of A with eigenvalue α , so $|\psi(\tau)\rangle = U|\psi(0)\rangle = e^{-i\alpha}|\psi(0)\rangle$. Since we are already given that $\psi(\tau) = e^{i\phi}\psi(0)$, it follows that $\phi \equiv -\alpha \pmod{2\pi}$; thus $\beta \equiv 0 \pmod{2\pi}$. Any desired multiple of 2π can always be selected for β by appropriate choice of the function $\tilde{\psi}(t)$. So, β does not contain any information at all about the specific evolution $\psi(t)$, and thus it is not a physically meaningful quantity.

It it interesting to note that the A&A paper [26] never actually shows that their quantity β can ever be different from 0 (mod 2π), although they do prove that β has some other "interesting" properties (such as being independent of the gauge of

the original trajectory) which of course are true trivially if β is always congruent to zero.

Thus, it seems that one implication of our results (assuming they are correct) is that Aharonov and Anandan's particular version (at least) of the "geometric phase" is a chimera, and does not really exist. Further study is needed to verify this conclusion more rigorously, and also to determine whether other definitions of the Berry phase might escape from it, and retain a useful physical meaning that relates in some way to our quantity α . Since many researchers have reported the experimental detection of Berry-type phases (*e.g.*, see [32]), it seems highly unlikely that our results will turn out to nullify all versions of the geometric phase for all quantum evolutions. However, as of this writing, the correct resolution of the apparent discrepancy between theory and experiment on this question is not yet clear.

14 Conclusion

In this paper, we have shown that any continuous trajectory of a normalized state vector can be measured by a real-valued quantity which we call the *effort* \mathcal{F} , which is given by the line integral, along the trajectory, of the imaginary component of the inner product between adjacent states along the trajectory. This quantity is basis-independent, and is numerically equal to the probability-weighted average phase angle accumulated by the basis state coefficients (in radians), and to twice the area swept out by the coefficients in the complex plane, and also to the action of the time-dependent Hamiltonian along the trajectory, in units of \hbar . This notion of effort can be easily extended to apply also to transformation trajectories U'(t) over time, as well as to an overall resulting unitary transform U, where it measures the difficulty \mathcal{D} or minimum effort (over available trajectories) required to implement the desired transform in the worst case (maximizing over the possible initial states). Our framework can be used to easily rederive a variety of related results obtained by earlier papers for various more specialized cases.

The major implication of these results is that there is indeed a very definite sense in which we can say that the physical concept of energy does indeed precisely correspond to the computational concept of the rate of computation, that is, we can validly say that energy *is* the rate of physical computing activity, defined as the rate of change of the state vector, according to the measure that we have described in this paper. Furthermore, we can validly say that physical action *is* (an amount of) computation, defined as the total amount of change of the state vector, in the sense we have defined.

What about different specific types of energy, and specific types of action? Later papers along this line of research will survey how different types of energy and action can validly be identified with computational activity that is engaged in different types of processes. For example, heat may be identified with energy whose detailed configuration information is unknown (is entropy), rest mass-energy can be identified with energy that is engaged in updating a system's internal state in its rest frame, potential energy with phase rotation due to emission/absorption of virtual particles, and so forth. As a preview, it turns out that we

can even make our computational interpretation consistent with special relativity by subdividing the energy of a moving body (in a given observer frame) into the *functional* energy Φ that is associated with updating the body's internal state (this turns out to be just the negative Lagrangian -L = H - pv) and a *motional* part M = pv (related to but not quite the same as kinetic energy) that is associated with conveying the body through space; relativistic momentum then turns out to be the motional computational effort exerted per unit distance traversed. Future papers will elaborate on these related themes in more depth.

It is hoped that the long-term outcome of this line of thought will be to eventually show how *all* physical concepts and quantities can be rigorously understood in a well-defined mathematical framework that is also simultaneously well-suited for describing physical implementations of desired computational processes. That is, we seek an eventual unifying mathematical foundation that is appropriate for not only physical science, but also for device-level computer engineering and for physics-based computer science. We expect that such a unifying perspective should greatly facilitate the future design and development of maximally efficient computers constructed from nanoscale (and perhaps, someday, even smaller) components, machines that attempt to harness the underlying computational resources provided by physics in the most efficient possible fashion.

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